

Enriched categories and their centers

Zhi-Hao Zhang

Wu Wen-Tsun Key Laboratory of Mathematics of Chinese Academy of Sciences,
School of Mathematical Sciences, University of Science and Technology of China

Shenzhen Institute for Quantum Science and Engineering,
Southern University of Science and Technology

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Enriched monoidal categories I: centers

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An **enriched category** ${}^{\mathcal{A}}\mathcal{L}$ consists of the following data:

- A monoidal category \mathcal{A} , called the base category or the **background category**;
- a set $\text{Ob}({}^{\mathcal{A}}\mathcal{L})$, whose elements are called objects;
- an object ${}^{\mathcal{A}}\mathcal{L}(x, y) \in \mathcal{A}$ for each $x, y \in \text{Ob}({}^{\mathcal{A}}\mathcal{L})$, called the hom space;
- a morphism $\circ: {}^{\mathcal{A}}\mathcal{L}(y, z) \otimes {}^{\mathcal{A}}\mathcal{L}(x, y) \rightarrow {}^{\mathcal{A}}\mathcal{L}(x, z)$ in \mathcal{A} for each $x, y, z \in \text{Ob}({}^{\mathcal{A}}\mathcal{L})$, called the composition;
- a morphism $1_x: \mathbb{1} = \mathbb{1}_{\mathcal{A}} \rightarrow {}^{\mathcal{A}}\mathcal{L}(x, x)$ for each $x \in \text{Ob}({}^{\mathcal{A}}\mathcal{L})$, called the identity;

and the composition is associative and unital.

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- Let \mathbf{Set} be the monoidal category of sets. Then a \mathbf{Set} -enriched category is an ordinary category.

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- Given a 1d quantum liquid phase (gapped/gapless topological phase with/without symmetries), its observables in the long wave length limit form an enriched category, called the **topological skeleton** of the phase. [[Kong-Zheng: A mathematical theory of gapless edges of 2d topological orders, 2019](#)]

An important class of examples are given by the so-called **canonical construction** [Kelly: *Adjunction for enriched categories*, 1969].

Let \mathcal{B} be a monoidal category and \mathcal{M} be a left \mathcal{B} -module. We say \mathcal{M} is **enriched** in \mathcal{B} if the right adjoint of $- \odot x: \mathcal{B} \rightarrow \mathcal{M}$ exists for each $x \in \mathcal{M}$, denoted by $[x, -]: \mathcal{M} \rightarrow \mathcal{B}$. The object $[x, y] \in \mathcal{B}$ is called the **internal hom** of x and y .

By the adjunctions $(- \odot x) \dashv [x, -]$ we have:

- for each $x \in \mathcal{M}$, there is a morphism $1_x: \mathbb{1} \rightarrow [x, x]$ in \mathcal{B} induced by the identity morphism $1_x: \mathbb{1} \odot x = x \rightarrow x$ in \mathcal{M} ;
- for every $x, y, z \in \mathcal{M}$, there is a morphism $\circ: [y, z] \otimes [x, y] \rightarrow [x, z]$ induced by

$$([y, z] \otimes [x, y]) \odot x \simeq [y, z] \odot ([x, y] \odot x) \rightarrow [y, z] \odot y \rightarrow z.$$

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Then there is an enriched category denoted by ${}^{\mathcal{B}}\mathcal{M}$, where:

- the set of objects is the same as \mathcal{M} , i.e., $\text{Ob}({}^{\mathcal{B}}\mathcal{M}) := \text{Ob}(\mathcal{M})$;
- the hom spaces are given by internal homs, i.e., ${}^{\mathcal{B}}\mathcal{M}(x, y) := [x, y] \in \mathcal{B}$;
- the composition and identity morphisms are given by the above constructions.

Another way to construct enriched categories is given by the operation called the **changing of the background category**:

- Recall that a lax-monoidal functor $F: \mathcal{A} \rightarrow \mathcal{B}$ between two monoidal categories is a functor $F: \mathcal{A} \rightarrow \mathcal{B}$ equipped with a natural transformation $F^2: F(-) \otimes F(-) \rightarrow F(- \otimes -)$ and a morphism $F^0: \mathbb{1}_{\mathcal{B}} \rightarrow F(\mathbb{1}_{\mathcal{A}})$ satisfying the usual conditions for a monoidal functor.
- Suppose ${}^{\mathcal{A}}\mathcal{L}$ is an \mathcal{A} -enriched category and $F: \mathcal{A} \rightarrow \mathcal{B}$ is a lax-monoidal functor. Then F induces a \mathcal{B} -enriched category ${}^{\mathcal{B}}\mathcal{L}$ with the same objects as ${}^{\mathcal{A}}\mathcal{L}$ and the hom spaces ${}^{\mathcal{B}}\mathcal{L}(x, y) := F({}^{\mathcal{A}}\mathcal{L}(x, y))$.
- In particular, when $\mathcal{B} = \text{Set}$ and $F = \mathcal{A}(\mathbb{1}, -)$, the ordinary category $\mathcal{L} := {}^{\text{Set}}\mathcal{L}$ is called the **underlying category** of ${}^{\mathcal{A}}\mathcal{L}$.

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For enriched categories, changing the background should also be viewed as a ‘functor’. So we have found a new definition of an enriched functor.

An **enriched functor** $|F: {}^{\mathcal{A}}\mathcal{L} \rightarrow {}^{\mathcal{B}}\mathcal{M}$ consists of the following data:

- a lax-monoidal functor $\hat{F}: \mathcal{A} \rightarrow \mathcal{B}$, called the **background changing functor**;
- a map $F: \text{Ob}({}^{\mathcal{A}}\mathcal{L}) \rightarrow \text{Ob}({}^{\mathcal{B}}\mathcal{M})$;
- a morphism $|F_{x,y}: \hat{F}({}^{\mathcal{A}}\mathcal{L}(x,y)) \rightarrow {}^{\mathcal{B}}\mathcal{M}(F(x), F(y))$ in \mathcal{B} for each $x, y \in {}^{\mathcal{A}}\mathcal{L}$;

such that $|F$ preserves the composition and identity morphisms.

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In particular, if $\mathcal{B} = \mathcal{A}$ and $\hat{F} = \text{id}_{\mathcal{A}}$, such an enriched functor $|F: {}^{\mathcal{A}}\mathcal{L} \rightarrow {}^{\mathcal{A}}\mathcal{M}$ is called an \mathcal{A} -functor. This is the traditional definition of an enriched functor (for example, see [\[Kelly: Basic concepts of enriched category theory, 1982\]](#)).

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Remark: Every enriched functor $|F: {}^{\mathcal{A}}\mathcal{L} \rightarrow {}^{\mathcal{B}}\mathcal{M}$ can be decomposed as $|F: {}^{\mathcal{A}}\mathcal{L} \rightarrow {}^{\mathcal{B}}\mathcal{L} \rightarrow {}^{\mathcal{B}}\mathcal{M}$, where the first functor only changes the background, and the second is a \mathcal{B} -functor.

The **underlying functor** $F: \mathcal{L} \rightarrow \mathcal{M}$ of an enriched functor $|F: {}^{\mathcal{A}}\mathcal{L} \rightarrow {}^{\mathcal{B}}\mathcal{M}$ is defined by the map $F: \text{Ob}(\mathcal{L}) \rightarrow \text{Ob}(\mathcal{M})$ and the map

$$F_{x,y}: \mathcal{L}(x,y) = \mathcal{A}(\mathbb{1}, {}^{\mathcal{A}}\mathcal{L}(x,y)) \xrightarrow{\hat{F}} \mathcal{B}(\hat{F}(\mathbb{1}), \hat{F}({}^{\mathcal{A}}\mathcal{L}(x,y))) \\ \xrightarrow{\mathcal{B}(\hat{F}^0, |F_{x,y})} \mathcal{B}(\mathbb{1}, {}^{\mathcal{B}}\mathcal{M}(F(x), F(y))) = \mathcal{M}(F(x), F(y)).$$

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Similarly, an **enriched natural transformation** is defined by a background changing natural transformation, which is monoidal, and an underlying natural transformation.

The enriched categories, enriched functors and enriched natural transformations form a 2-category ${}^{\text{lax}}\mathbf{ECat}$. It is a symmetric monoidal 2-category with the Cartesian product defined by

- $\text{Ob}({}^{\mathcal{A}}\mathcal{L} \times {}^{\mathcal{B}}\mathcal{M}) := \text{Ob}({}^{\mathcal{A}}\mathcal{L}) \times \text{Ob}({}^{\mathcal{B}}\mathcal{M})$;
- $({}^{\mathcal{A}}\mathcal{L} \times {}^{\mathcal{B}}\mathcal{M})((x, y), (x', y')) := ({}^{\mathcal{A}}\mathcal{L}(x, x'), {}^{\mathcal{B}}\mathcal{M}(y, y')) \in \mathcal{A} \times \mathcal{B}$.

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Remark: Fix a monoidal category \mathcal{A} . Then there is a 2-category ${}^{\mathcal{A}}\mathbf{Cat}$ of \mathcal{A} -enriched categories, \mathcal{A} -functors and \mathcal{A} -natural transformations. This is the traditional definition of a 2-category of enriched categories (for example, see [[Kelly: Basic concepts of enriched category theory, 1982](#)]).

We define an **enriched monoidal category**, an **enriched braided monoidal category** and an **enriched symmetric monoidal category** as an E_1 -algebra, an E_2 -algebra and an E_3 -algebra in the symmetric monoidal 2-category **ECat**.

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Intuitions for E_n -algebras

- We can imagine an E_n -algebra as a collection of particles living in an n -dimensional space. So these particles can be fused together in n directions.
- Intuitively, an E_n -algebra in a symmetric monoidal ∞ -category \mathcal{C} is an object equipped with n compatible multiplications.
- In particular, an E_0 -algebra in \mathcal{C} is an object $x \in \mathcal{C}$ equipped with a morphism $\mathbb{1} \rightarrow x$.

Examples:

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- (b) In the symmetric monoidal 2-category **Cat** of categories, functors and natural transformations, an E_1 -algebra is a monoidal category, an E_2 -algebra is a braided monoidal category and an E_n -algebra for $n \geq 3$ is a symmetric monoidal category.

Taking the underlying category $\mathcal{A}|\mathcal{L} \mapsto \mathcal{L}$ is a symmetric monoidal 2-functor $\mathbf{ECat} \mapsto \mathbf{Cat}$, thus it maps an E_n -algebra to an E_n -algebra. Similarly, taking the background category $\mathcal{A}|\mathcal{L} \mapsto \mathcal{A}$ is also a symmetric monoidal 2-functor $\mathbf{ECat} \mapsto \mathbf{Alg}_{E_1}(\mathbf{Cat})$.

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- (a) Given an enriched monoidal category $\mathcal{A}|\mathcal{L}$, its underlying category \mathcal{L} is a monoidal category, and its background category \mathcal{A} is a braided monoidal category.
- (b) Given an enriched braided monoidal category $\mathcal{A}|\mathcal{L}$, its underlying category \mathcal{L} is a braided monoidal category, and its background category \mathcal{A} is a symmetric monoidal category.
- (c) Given an enriched symmetric monoidal category $\mathcal{A}|\mathcal{L}$, its underlying category \mathcal{L} and background category \mathcal{A} are both symmetric monoidal categories.

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2. Let \mathcal{C} be a braided monoidal category. A **monoidal left $\bar{\mathcal{C}}$ -module** is a monoidal category \mathcal{M} equipped with a braided monoidal functor $\varphi: \bar{\mathcal{C}} \rightarrow \mathfrak{Z}_1(\mathcal{M})$. If \mathcal{M} is enriched in $\bar{\mathcal{C}}$, the canonical construction ${}^{\mathcal{C}}\mathcal{M}$ is an enriched monoidal category.

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3. Let \mathcal{E} be a symmetric monoidal category. A **braided left \mathcal{E} -module** is a braided monoidal category \mathcal{M} equipped with a symmetric monoidal functor $\phi: \mathcal{E} \rightarrow \mathfrak{Z}_2(\mathcal{M})$. If \mathcal{M} is enriched in \mathcal{E} , the canonical construction ${}^{\mathcal{E}}\mathcal{M}$ is an enriched braided monoidal category. Moreover, ${}^{\mathcal{E}}\mathcal{M}$ is symmetric if \mathcal{M} is symmetric.

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Example: centers in the symmetric monoidal 1-category Vec

- (a) The center of a vector space V is given by the endomorphism algebra $\text{End}(V)$: a linear map $X \otimes V \rightarrow V$ is equivalent to a linear map $X \rightarrow \text{Hom}(V, V) = \text{End}(V)$.

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- (b) In Vec , the E_1 -center of a \mathbb{k} -algebra A is the usual center

$$Z(A) := \{z \in A \mid za = az, \forall a \in A\}.$$

Indeed, if $f: X \otimes A \rightarrow A$ is both an algebra homomorphism and a left unital action, then

$$f(x \otimes 1)a = f(x \otimes 1)f(1 \otimes a) = f(x \otimes a) = f(1 \otimes a)f(x \otimes 1) = af(x \otimes 1).$$

Example: centers in the symmetric monoidal 2-category **Cat**:

- (a) The center of a category \mathcal{X} is given by the endofunctor category $\text{End}(\mathcal{X})$.
- (b) The E_1 -center of a monoidal category \mathcal{A} is given by the Drinfeld center (or monoidal center) $\mathfrak{Z}_1(\mathcal{A})$, whose objects are objects in \mathcal{A} equipped with a half-braiding.
- (c) The E_2 -center of a braided monoidal category \mathcal{C} is given by the Müger center (or symmetric center)

$$\mathfrak{Z}_2(\mathcal{C}) := \{x \in \mathcal{C} \mid c_{y,x} \circ c_{x,y} = \text{id}_{x \otimes y}, \forall y \in \mathcal{C}\}.$$

Here we give the centers of enriched (monoidal) categories obtained from the canonical construction:

1. Suppose \mathcal{B} is a rigid monoidal category and \mathcal{M} is a left \mathcal{B} -module. Then $\text{Fun}_{\mathcal{B}}(\mathcal{M}, \mathcal{M})$ is a left $\mathfrak{Z}_1(\mathcal{B})$ -module. Moreover, $\text{Fun}_{\mathcal{B}}(\mathcal{M}, \mathcal{M})$ is a left monoidal $\overline{\mathfrak{Z}_1(\mathcal{B})}$ -module, i.e., there is a braided monoidal functor $\overline{\mathfrak{Z}_1(\mathcal{B})} \rightarrow \mathfrak{Z}_1(\text{Fun}_{\mathcal{B}}(\mathcal{M}, \mathcal{M}))$. The center of ${}^{\mathcal{B}}\mathcal{M}$ is given by the enriched monoidal category ${}^{\mathfrak{Z}_1(\mathcal{B})}\text{Fun}_{\mathcal{B}}(\mathcal{M}, \mathcal{M})$.

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2. Suppose \mathcal{C} is a braided monoidal category and \mathcal{M} is a monoidal left $\overline{\mathcal{C}}$ -module defined by a braided monoidal functor $\varphi: \overline{\mathcal{C}} \rightarrow \mathfrak{Z}_1(\mathcal{M})$. Then the Müger centralizer $\mathfrak{Z}_2(\varphi)$ of \mathcal{C} in $\mathfrak{Z}_1(\mathcal{M})$ is a braided left $\mathfrak{Z}_2(\mathcal{C})$ -module, i.e., there is a symmetric monoidal functor $\mathfrak{Z}_2(\mathcal{C}) \rightarrow \mathfrak{Z}_2(\mathfrak{Z}_2(\varphi))$. The E_1 -center of ${}^{\mathcal{C}}\mathcal{M}$ is ${}^{\mathfrak{Z}_2(\mathcal{C})}\mathfrak{Z}_2(\varphi)$.

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3. The E_2 -center of an enriched braided monoidal category ${}^{\mathcal{A}}\mathcal{L}$ is given by ${}^{\mathcal{A}}\mathfrak{Z}_2(\mathcal{L})$.

In particular, we have the following results.

- (a) Let \mathcal{A} be an indecomposable multi-fusion category and \mathcal{L} be a finite semisimple left \mathcal{A} -module. Then

$$\mathfrak{Z}_1(\mathfrak{Z}_0({}^{\mathcal{A}}\mathcal{L})) \simeq \mathfrak{Z}_1({}^{\mathfrak{Z}_1(\mathcal{A})}\text{Fun}_{\mathcal{A}}(\mathcal{L}, \mathcal{L})) \simeq {}^{\text{Vec}}\text{Vec} = \text{Vec}.$$

- (b) Let \mathcal{C} be a nondegenerate braided fusion category and \mathcal{M} be an indecomposable multi-fusion left $\overline{\mathcal{C}}$ -module defined by a braided tensor functor $\varphi: \overline{\mathcal{C}} \rightarrow \mathfrak{Z}_1(\mathcal{M})$.
Then

$$\mathfrak{Z}_2(\mathfrak{Z}_1({}^{\mathcal{C}}\mathcal{M})) \simeq \mathfrak{Z}_2({}^{\text{Vec}}\mathfrak{Z}_2(\varphi)) \simeq {}^{\text{Vec}}\text{Vec} = \text{Vec}.$$

The center of a center is trivial. Its physical meaning is that the bulk of a bulk is trivial.

Thanks!