

Sustech-Nagoya workshop on Quantum Science

A 2-Categorical Interpretation of Reconstruction Theorem for Weak Hopf Algebras

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based on joint work with Zhi-Hao Zhang (in preparation).

The notion of weak Hopf algebras (WHAs) originates from the motivation to describe certain symmetries in quantum field theory and operator algebras [Böhm-Nill-Szlachányi:99][Mack&Schomerus:92][Schomerus:95]. As a “quantum symmetry”, WHAs along with their representations have wide applications in the studies of both quantum phenomena[Kitaev-Kong:1104.5047] and novel algebraic structures[Nikshych&Vainerman:00].

Our goal today is to give a new interpretation/generalisation of a basic theorem on WHAs. This theorem is called the *reconstruction theorem* (or Tannaka-Krein duality) for WHAs, of which the original version is due to Prof. Takahiro Hayashi[Hayashi:math/9904073][Ostrik:math/0111139], who is also from Nagoya University. This theorem establishes a strong relation between WHAs and their representations.

- Basics on WHAs.
- The reconstruction theorem.
- An interpretation.

Basics on WHAs

(Finite dim'l) Weak Hopf algebras $H(M, u, \Delta, \varepsilon, S)$

$$M: H \otimes H \rightarrow H \quad u: \mathbb{C} \rightarrow H$$

$$\Delta: H \rightarrow H \otimes H \quad \varepsilon: H \rightarrow \mathbb{C} \quad S: H \rightarrow H$$

- (H, M, u) is a (finite dimensional) associative algebra.
- (H, Δ, ε) is a coassociative coalgebra:

$$(\Delta \otimes \text{id})\Delta = (\text{id} \otimes \Delta)\Delta \quad (\varepsilon \otimes \text{id})\Delta = \text{id} = (\text{id} \otimes \varepsilon)\Delta.$$

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- Δ preserves multiplication (where $\tau: H \otimes H \rightarrow H \otimes H, a \otimes b \mapsto b \otimes a$):

$$\Delta M = (M \otimes M)(\text{id} \otimes \tau \otimes \text{id})(\Delta \otimes \Delta).$$

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- M, u, Δ and ε satisfy:

$$(\Delta \otimes \text{id})\Delta u = (\text{id} \otimes M \otimes \text{id})(\Delta \otimes \Delta)(u \otimes u) = (\text{id} \otimes (M \circ \tau) \otimes \text{id})(\Delta \otimes \Delta)(u \otimes u)$$

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- The antipode S is an algebra-antihomomorphism and also a coalgebra-antihomomorphism, subject to certain compatibility conditions.

A simple example: groupoid algebra

\mathcal{G} : a finite groupoid. Define a weak Hopf algebra $\mathbb{C}[\mathcal{G}] := \text{span}\{g \mid g \in \text{Mor}(\mathcal{G})\}$ with

$$M(g_1 \otimes g_2) = \begin{cases} g_1 g_2 & \text{if } g_1 \text{ can be left composed to } g_2; \\ 0 & \text{otherwise.} \end{cases}$$

$$u(1) = \sum_{a \in \text{Ob}(\mathcal{G})} \text{id}_a$$

$$\Delta(g) = g \otimes g$$

$$\varepsilon(g) = 1$$

$$S(g) = g^{-1}.$$

Remark. When \mathcal{G} has only one object, i.e. is a group, then $\mathbb{C}[\mathcal{G}]$ is the group algebra and is a basic example of *Hopf algebras*.

The two canonical subalgebras

An interesting fact is that there are two canonical subalgebras of a WHA H . To see this, first note that there are two projections

$$p_L: H \rightarrow H, x \mapsto \varepsilon(x1_1)1_2$$

and

$$p_R: H \rightarrow H, x \mapsto 1_1\varepsilon(1_2x),$$

where we have used Sweedler's notation $\Delta(1) \equiv 1_1 \otimes 1_2$. Then we have the following facts:

1. The images $H_L := p_L(H)$ and $H_R := p_R(H)$ are closed under multiplication and contains $1 \in H$, thus are **subalgebras** of H . Moreover, they are **separable**.
2. The algebras H_L and H_R are anti-isomorphic, i.e., $H_R \cong H_L^{\text{op}}$.
3. The algebras H_L and H_R **mutually commute** with each other, i.e., for $x \in H_L$ and $y \in H_R$, we have $xy = yx$.

The two canonical subalgebras (cont'd)

4. There is $S(H_L) = H_R, S(H_R) = H_L$, i.e., the antipode “transfers” the two subalgebras to each other.

The case: groupoid algebra

Consider the groupoid algebra $\mathbb{C}[\mathcal{G}]$. For $g: a \rightarrow b \in \text{Mor}(\mathcal{G})$, the two projections read

$$\begin{aligned} p_L: g &\mapsto \text{id}_a \\ p_R: g &\mapsto \text{id}_b, \end{aligned}$$

and consequently

$$H_L = H_R = \text{span}\{\text{id}_a \mid a \in \text{Ob}(\mathcal{G})\} = \mathbb{C}^{\oplus |\text{Ob}(\mathcal{G})|}.$$

The projections p_L and p_R are hence called the **source** and **target** maps respectively, while subalgebras H_L and H_R the **source** and **target** algebras.

Remark. In this case the source and target algebras are commutative algebras, however this is *not* true in a generic WHA.

**The representation category of a
WHA and the reconstruction
theorem**

The representation category of H

Let H be a WHA and let ${}_H\text{Mod}$ denote the category of left modules over H . Then the aforementioned facts on the two subalgebras lead to the following statement:

Fact

There is a forgetful functor $U: {}_H\text{Mod} \rightarrow {}_{H_L \otimes H_R}\text{Mod} = {}_{H_L}\text{Mod}_{H_L}$.

Here the functor U acts in the following way: given a left H -module V , the left $H_L \otimes H_R$ -module $U(V)$ has commuting left actions $(x, v) \mapsto x.v$ and $(y, v) \mapsto y.v$ for $x \in H_L, y \in H_R$.

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Note that ${}_{H_L}\text{Mod}_{H_L}$ is a monoidal category whose tensor product is given by the relative tensor product over H_L . Using axioms of WHA, one can also show that

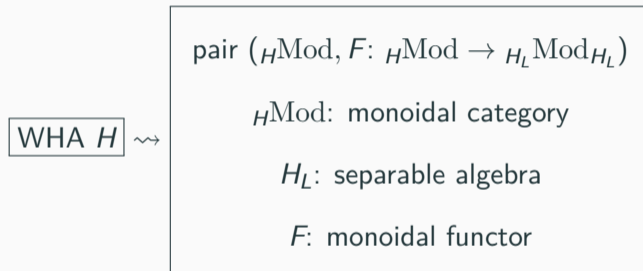
Fact (non-trivial)

(1) The category ${}_H\text{Mod}$ carries a monoidal structure; (2) The functor U is monoidal.

The reconstruction theorem

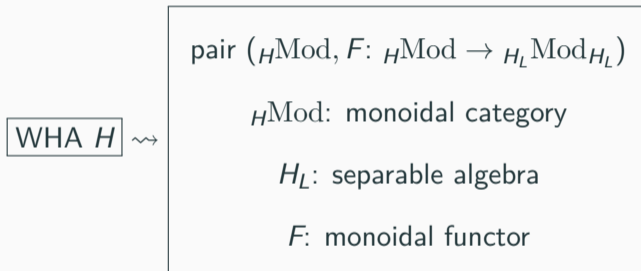
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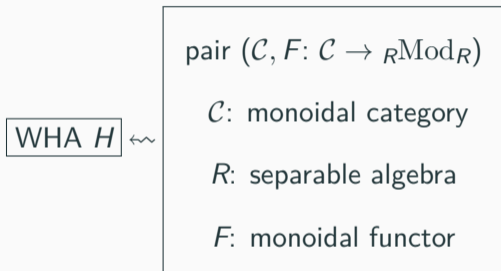


Theorem (Reconstruction theorem for WHAs)

The “ \rightsquigarrow ” has an inverse.

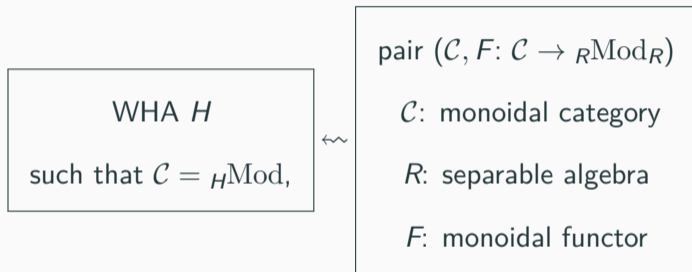
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What should the inverse of “ \rightsquigarrow ” achieve? The inverse of “ \rightsquigarrow ” should read



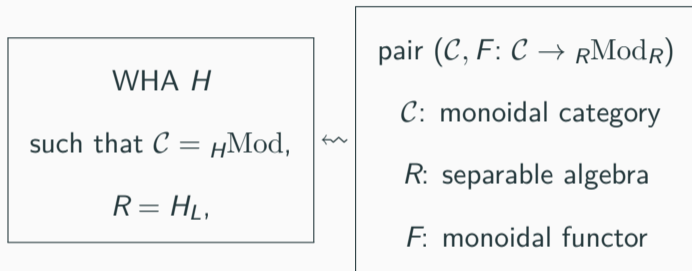
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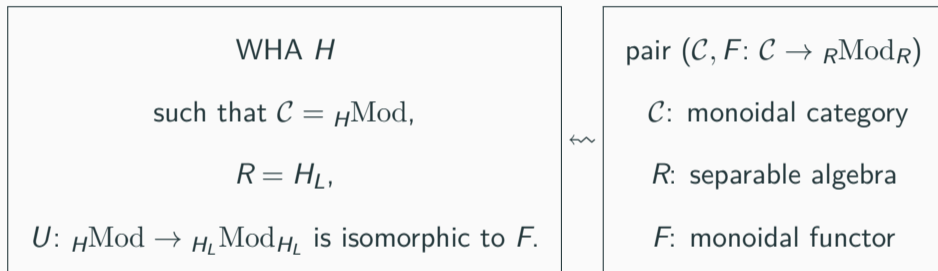
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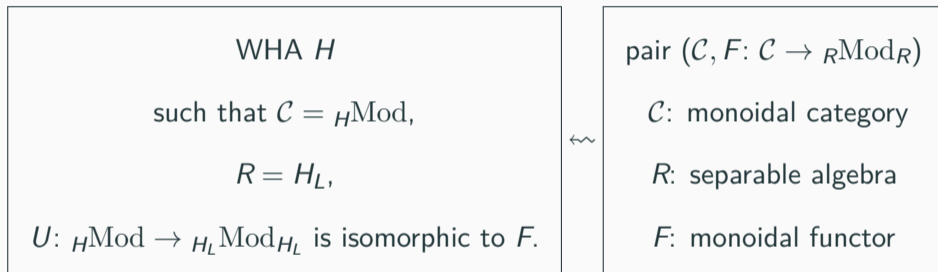
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Now I briefly show you how “ $\leftarrow \rightsquigarrow$ ” can be realized, i.e., “proof” the theorem.

The reconstruction theorem (cont'd)

Define a long forgetful functor:

$$\tilde{F}: \mathcal{C} \xrightarrow{F} {}_R\text{Mod}_R \xrightarrow{f} \text{Vec},$$

where Vec denotes the category of \mathbb{C} -vector spaces and f forgets the R -bimodule action. Then there is a standard construction $H := \text{End}(\tilde{F})$ such that ${}_H\text{Mod} = \mathcal{C}$ as categories, where $\text{End}(\tilde{F})$ denotes the algebra of endo-natural transformations of \tilde{F} .

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It still remains to give the weak Hopf algebra structure on H . However, I will skip most of it, and only guide you to see that there are two subalgebras of H isomorphic to R and R^{op} respectively. For $r \in R$, define natural transformations $r^\sharp, r^\flat: \tilde{F} \Rightarrow \tilde{F}$ as follows:

$$\begin{aligned} (r^\sharp)_X: \tilde{F}(X) &\rightarrow \tilde{F}(X) & (r^\flat)_X: \tilde{F}(X) &\rightarrow \tilde{F}(X) \\ v &\mapsto r \cdot v & v &\mapsto v \cdot r, \end{aligned}$$

where $X \in \mathcal{C}$, and the (left and right) actions of r on v is defined in the bimodule $F(X)$.

The reconstruction theorem (cont'd)

Define sets

$$R^\sharp := \{r^\sharp \mid r \in R\}, \quad R^\flat := \{r^\flat \mid r \in R\},$$

then it is not hard to see that $R^\sharp, R^\flat \subset \text{End}(\tilde{F}) = H$ are subalgebras isomorphic to R, R^{op} respectively.

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then it is not hard to see that $R^\sharp, R^\flat \subset \text{End}(\tilde{F}) = H$ are subalgebras isomorphic to R, R^{op} respectively. Not surprisingly, if we work out all structures of H , which allows us to compute H_L and H_R , then $H_L = R^\sharp$ and $H_R = R^\flat$. Thus we have

$$\boxed{U: {}_H\text{Mod} \rightarrow {}_{H_L}\text{Mod}_{H_L}} = \boxed{F: \mathcal{C} \rightarrow {}_R\text{Mod}_R}.$$

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An interpretation

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I want to emphasize that this reconstruction theorem is essentially a 2-categorical phenomenon. Define the following 2-category 2Vec :

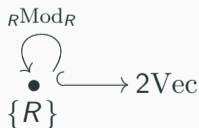
1. 0-,1-,2-morphisms are separable algebras, bimodules and bimodule maps respectively. More precisely, we define the 1-hom-category from R to R' to be ${}_{R'}\text{Mod}_R$ for separable algebras R, R' .
2. Composition of 1-morphisms are given by relative tensor product. More precisely, the composition functor ${}_{R'}\text{Mod}_{R'} \times {}_{R'}\text{Mod}_R \rightarrow {}_{R'}\text{Mod}_R$ is defined to be $\otimes_{R'}$.

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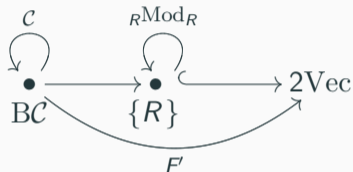
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Then it is easy to see that the monoidal category ${}_R\text{Mod}_R$ can be viewed as the looping of the one-object full sub-2-category $\{R\}$ of $2Vec$, as is shown in the graph below.



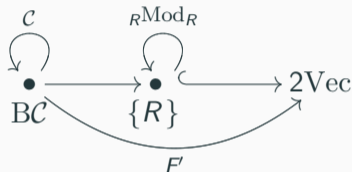
An interpretation (cont'd)

In particular, the data of the pair $(\mathcal{C}, F: \mathcal{C} \rightarrow {}_R\text{Mod}_R)$ are precisely the data of a pair $(BC, F': BC \rightarrow 2\text{Vec})$, where BC and F' are defined in the following figure.



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Thus it is equivalent to say that the reconstruction starts really from the pair (BC, F') . A natural question is: **can one reconstruct some algebraic data from more general 2-functors to 2Vec ?**

Remark. Why do we care about 2Vec ? The 2-category 2Vec is 2-equivalent to the category of finite semisimple \mathbb{C} -linear categories, \mathbb{C} -linear functors and natural transformations, which deserves to be viewed as the categorification of Vec .

An interpretation (cont'd)

The answer is yes. Let \mathcal{D} be a 2-category and let $U: \mathcal{D} \rightarrow 2\text{Vec}$ be a 2-functor. We denote the hom-categories in \mathcal{D} by $\{\mathcal{D}(a, b)\}_{a, b \in \text{Ob}(\mathcal{D})}$, and the local 1-functors of U by $\{U_{a, b}: \mathcal{D}(a, b) \rightarrow U_b \text{Mod}_{U_a}\}_{a, b \in \text{Ob}(\mathcal{D})}$. Then we have a 1-functor

$$\widetilde{U}_{a, b}: \mathcal{D}(a, b) \xrightarrow{U_{a, b}} U_b \text{Mod}_{U_a} \xrightarrow{f} \text{Vec}$$

for objects $a, b \in \mathcal{D}$, where f denotes the functor forgetting the bimodule actions.

Then the standard construction tells us that the algebra $H_{a, b} := \text{End}(\widetilde{U}_{a, b})$ satisfies that $H_{a, b} \text{Mod} = \mathcal{D}(a, b)$ as categories. Moreover, there are algebra homomorphisms $U_b \rightarrow H_{a, b}$ and $U_a^{\text{op}} \rightarrow H_{a, b}$. As a matter of fact, the whole algebraic data reconstructed from U can be formulated.

Definition (first half)

A *weak Hopf algebra* \mathcal{H} consists of the following data:

1. A set of objects $\text{Ob}(\mathcal{H})$;
2. For objects $a, b \in \text{Ob}(\mathcal{H})$, there is an algebra $H_{a,b}$;
3. For objects $a, b, c \in \text{Ob}(\mathcal{H})$, a *generalized comultiplication*
 $\Delta^{abc}: H_{a,c} \rightarrow H_{b,c} \otimes H_{a,b}$;
4. For object $a \in \text{Ob}(\mathcal{H})$, a counit $\varepsilon^a: H_{a,a} \rightarrow \mathbb{k}$;
5. For objects $a, b \in \text{Ob}(\mathcal{H})$, a *generalized antipode* $S_{a,b}: H_{a,b} \rightarrow H_{b,a}$,

which satisfy the following conditions for all $a, b, c, d \in \text{Ob}(\mathcal{H})$

- Coassociativity and counitality hold (hence $\{H_{a,b}\}_{a,b \in \text{Ob}(\mathcal{H})}$ forms a *cocategory*):

$$(\Delta^{bcd} \otimes \text{id}) \circ \Delta^{abd} = (\text{id} \otimes \Delta^{abc}) \circ \Delta^{acd}, \quad (\varepsilon^b \otimes \text{id}) \circ \Delta^{abb} = \text{id} = (\text{id} \otimes \varepsilon^a) \circ \Delta^{aab};$$

Definition (second half)

- The generalized comultiplication $\Delta^{abc} : H_{a,c} \rightarrow H_{b,c} \otimes H_{a,b}$ preserves multiplication:

$$\Delta^{abc} M_{a,c} = (M_{b,c} \otimes M_{a,b})(\text{id} \otimes \tau \otimes \text{id})(\Delta^{abc} \otimes \Delta^{abc}),$$

where $M_{a,c}$ denotes the multiplication of $H_{a,c}$, etc.;

- The generalized comultiplications and counits satisfy other compatibilities:

$$\begin{aligned} (\text{id} \otimes M_{b,c} \otimes \text{id})(\Delta^{bcd} \otimes \Delta^{abc})(u_{b,d} \otimes u_{a,c}) &= (\Delta^{bcd} \otimes \text{id})\Delta^{abd} u_{a,d} = \\ (\text{id} \otimes (M_{b,c}\tau) \otimes \text{id})(\Delta^{bcd} \otimes \Delta^{abc})(u_{b,d} \otimes u_{a,c}) & \\ (M_{a,b} \otimes \varepsilon^a M_{a,a})(\text{id} \otimes \Delta^{aab} u_{a,b} \otimes \text{id}) &= (\text{id} \otimes \varepsilon^a M_{a,a})(\Delta^{aab} \otimes \text{id}) \\ (\varepsilon^a M_{a,a} \otimes M_{a,b})(\text{id} \otimes (\tau \Delta^{aab} u_{a,b}) \otimes \text{id}) &= (\text{id} \otimes \varepsilon^a M_{a,a})(\tau \otimes \text{id})(\text{id} \otimes \Delta^{aab}) \\ (\varepsilon^a M_{a,a} \otimes M_{c,a})(\text{id} \otimes \Delta^{caa} u_{c,a} \otimes \text{id}) &= (\varepsilon^a M_{a,a} \otimes \text{id})(\text{id} \otimes \Delta^{caa}) \\ (\text{id} \otimes \varepsilon^a M_{a,a})(\text{id} \otimes (\tau \Delta^{caa} u_{c,a}) \otimes \text{id}) &= (\varepsilon^a M_{a,a} \otimes \text{id})(\text{id} \otimes \tau)(\Delta^{caa} \otimes \text{id}); \end{aligned}$$

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- The antipode $S_{a,b} : H_{a,b} \rightarrow H_{b,a}$ is an algebra anti-homomorphism compatible with the comultiplications, and satisfies certain conditions.

Definition (second half)

- The generalized comultiplication $\Delta^{abc} : H_{a,c} \rightarrow H_{b,c} \otimes H_{a,b}$ preserves multiplication:

$$\Delta^{abc} M_{a,c} = (M_{b,c} \otimes M_{a,b})(\text{id} \otimes \tau \otimes \text{id})(\Delta^{abc} \otimes \Delta^{abc}),$$

where $M_{a,c}$ denotes the multiplication of $H_{a,c}$, etc.;

- The generalized comultiplications and counits satisfy other compatibilities:

$$\begin{aligned} (\text{id} \otimes M_{b,c} \otimes \text{id})(\Delta^{bcd} \otimes \Delta^{abc})(u_{b,d} \otimes u_{a,c}) &= (\Delta^{bcd} \otimes \text{id})\Delta^{abd} u_{a,d} = \\ (\text{id} \otimes (M_{b,c}\tau) \otimes \text{id})(\Delta^{bcd} \otimes \Delta^{abc})(u_{b,d} \otimes u_{a,c}) & \\ (M_{a,b} \otimes \varepsilon^a M_{a,a})(\text{id} \otimes \Delta^{aab} u_{a,b} \otimes \text{id}) &= (\text{id} \otimes \varepsilon^a M_{a,a})(\Delta^{aab} \otimes \text{id}) \\ (\varepsilon^a M_{a,a} \otimes M_{a,b})(\text{id} \otimes (\tau \Delta^{aab} u_{a,b}) \otimes \text{id}) &= (\text{id} \otimes \varepsilon^a M_{a,a})(\tau \otimes \text{id})(\text{id} \otimes \Delta^{aab}) \\ (\varepsilon^a M_{a,a} \otimes M_{c,a})(\text{id} \otimes \Delta^{caa} u_{c,a} \otimes \text{id}) &= (\varepsilon^a M_{a,a} \otimes \text{id})(\text{id} \otimes \Delta^{caa}) \\ (\text{id} \otimes \varepsilon^a M_{a,a})(\text{id} \otimes (\tau \Delta^{caa} u_{c,a}) \otimes \text{id}) &= (\varepsilon^a M_{a,a} \otimes \text{id})(\text{id} \otimes \tau)(\Delta^{caa} \otimes \text{id}); \end{aligned}$$

- The antipode $S_{a,b} : H_{a,b} \rightarrow H_{b,a}$ is an algebra anti-homomorphism compatible with the comultiplications, and satisfies certain conditions.

Remark. A one-object weak Hopf algebra is a weak Hopf algebra.

Theorem

One can reconstruct a weak Hopf algebroid from a pair $(\mathcal{D}, U: \mathcal{D} \rightarrow 2\text{Vec})$ where \mathcal{D} is a 2-category and U is a 2-functor.

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One can reconstruct a weak Hopf algebroid from a pair $(\mathcal{D}, U: \mathcal{D} \rightarrow 2\text{Vec})$ where \mathcal{D} is a 2-category and U is a 2-functor.

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Theorem (Part II)

Given a weak Hopf algebroid \mathcal{H} , there is a pair $(\text{Mod}(\mathcal{H}), \mathcal{F}_{\mathcal{H}})$ where $\text{Mod}(\mathcal{H})$ is a 2-category and $\mathcal{F}_{\mathcal{H}}: \text{Mod}(\mathcal{H}) \rightarrow 2\text{Vec}$ is a 2-functor.

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This theorem generalizes the reconstruction theorem for WHAs. More precisely, if we take \mathcal{D} to be a one-object 2-category, then we recover the latter. This provides a way for understanding the original theorem, especially the appearance of the category ${}_R\text{Mod}_R$ there.

Some remarks on weak Hopf algebroids

- Let \mathcal{H} be a weak Hopf algebroid. For any object $a \in \text{Ob}(\mathcal{H})$, the algebra $H_{a,a}$ is a weak Hopf algebra. We denote the canonical subalgebra $(H_{a,a})_L \subset H_{a,a}$ by H_a .
- The data of the algebra homomorphisms $H_b \rightarrow H_{a,b}$ and $H_a^{\text{op}} \rightarrow H_{a,b}$ can be obtained by structures of \mathcal{H} : they are the following maps restricted on H_b and H_a respectively

$$\begin{aligned} \rho_L: H_{b,b} &\rightarrow H_{a,b}, \quad x \mapsto \varepsilon^b(x 1_1^{ab}) 1_2^{ab} \\ \text{and } \rho_R: H_{a,a} &\rightarrow H_{a,b}, \quad x \mapsto 1_1^{ab} \varepsilon^a(1_2^{ab} x). \end{aligned}$$

- Some examples of weak Hopf algebras can have its weak Hopf algebroid counterpart.
- In particular, the weak Hopf algebra studied in Kitaev and Kong (2012) is indeed a substructure of a weak Hopf algebroid, which can be used to describe more general boundary excitations in Levin-Wen models [[Levin&Wen:cond-mat/0404617](#)].

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