

A Geometric Study of the Split Decomposition

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Abstract

This paper sheds a new light on the split decomposition theory and T-theory from the viewpoint of convex analysis and polyhedral geometry. By regarding finite metrics as discrete concave functions, Bandelt-Dress' split decomposition can be derived as a special case of more general decomposition of polyhedral/discrete concave functions introduced in this paper. It is shown that the combinatorics of splits discussed in connection to the split decomposition corresponds to the geometric properties of a hyperplane arrangement and a point configuration. Using our approach, the split decomposition of metrics can be naturally extended to distance functions, which may violate the triangle inequality, using partial split distances.

1 Introduction

Mathematical treatment of metric spaces on a finite set (*finite metric spaces*) has come to be increasingly important in the area of bioinformatics and phylogenetics; see [2],[24]. The central problem in phylogenetics is reconstructing phylogenetic trees from given experimental data, e.g., DNA sequences. If the data is given as a distance matrix expressing dissimilarity between species, the problem is to search for a tree metric that “fits” the given distance matrix, where a metric is called a *tree metric* if it can be represented as the path metric on some weighted tree.

T-theory [9], developed by A. Dress and coworkers, provides a beautiful mathematical framework for this phylogenetic problem. The *split decomposition*, due to Bandelt and Dress [1], is a T-theoretical phylogenetic analysis method, which is briefly summarized as follows; see Section 4 for precise definitions. Let V be a finite set and $d : V \times V \rightarrow \mathbf{R}$ a metric on V . Then the split decomposition decomposes d into the following form:

$$d = \sum_{\{A,B\}:\text{split}} \alpha_{\{A,B\}}^d \delta_{\{A,B\}} + d', \quad (1.1)$$

where a *split* means a bipartition of V , $\delta_{\{A,B\}}$ is the *split metric* defined in (4.3), $\alpha_{\{A,B\}}^d$ is the nonnegative number called *isolation index* defined in (4.4), and d' is a *split-prime metric*.

One of the main aims of this paper is to derive the split decomposition in a natural way as a special case of a decomposition of (discrete) convex/concave functions that we propose in this paper. This view point of convex analysis sheds a new light on such important concepts as split decomposition, tight span, isolation index, weakly compatible splits. Furthermore, our approach leads to a natural extension of the split decomposition for distances, which may violate the triangle inequality.

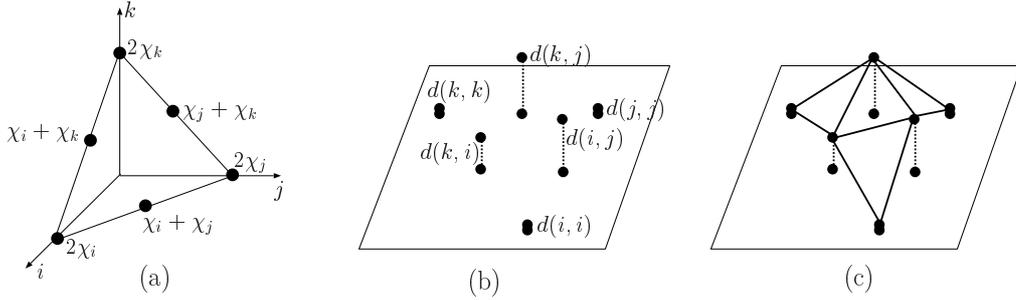


Figure 1: 3-dimensional example

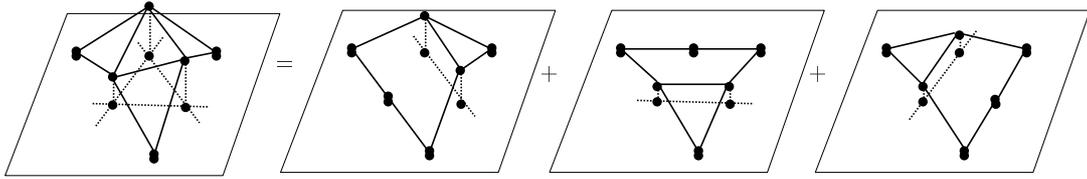


Figure 2: Decomposition of the concave function associated with d

To describe our key idea to understand the split decomposition from the view point of polyhedral geometry, let us take the following small example. Let $V = \{i, j, k\}$ be a 3-point set and $d : V \times V \rightarrow \mathbf{R}$ a metric on V defined as

$$d = \begin{array}{c|ccc} & i & j & k \\ \hline i & 0 & 4 & 2 \\ j & 4 & 0 & 4 \\ k & 2 & 4 & 0 \end{array} . \quad (1.2)$$

Consider the point set $\{(\chi_i + \chi_j, d(i, j)) \mid i, j \in V\} \subseteq \mathbf{R}^V \times \mathbf{R}$ and the upper envelope of its convex hull, where χ_i is the i th unit vector. Namely, we regard d as a function defined on the point set $\Lambda = \{\chi_i + \chi_j \mid i, j \in V\}$. We illustrate this situation in Figure 1. The point set Λ consists of vertices of the simplex and midpoints of its edges, and lies on the hyperplane $\{x \in \mathbf{R}^V \mid \sum_{i \in V} x(i) = 2\}$ (Figure 1 (a)). Hence, we can project $\{(\chi_i + \chi_j, d(i, j)) \mid i, j \in V\}$ to 3-dimensional space (Figure 1 (b)).

We observe that the concave function of this upper envelope can be decomposed into a sum of three simple concave functions as Figure 2. Then this decomposition of the concave function associated with d induces the decomposition of d as

$$\begin{aligned} d &= 2 \begin{pmatrix} 0 & 1 & 0 \\ 1 & 0 & 1 \\ 0 & 1 & 0 \end{pmatrix} + \begin{pmatrix} 0 & 1 & 1 \\ 1 & 0 & 0 \\ 1 & 0 & 0 \end{pmatrix} + \begin{pmatrix} 0 & 0 & 1 \\ 0 & 0 & 1 \\ 1 & 1 & 0 \end{pmatrix} \\ &= 2\delta_{\{2\},\{1,3\}} + \delta_{\{1\},\{2,3\}} + \delta_{\{3\},\{1,2\}} \end{aligned} \quad (1.3)$$

As will be shown in this paper, this decomposition coincides with Bandelt-Dress' split decomposition.

The observation above motivates us to study the following type of a decomposition of a polyhedral convex function f on \mathbf{R}^n as

$$f(x) = \sum_{(a,r) \in \mathbf{R}^n \times \mathbf{R}} c_{a,r}^f |\langle a, x \rangle - r| + f'(x) \quad (x \in \mathbf{R}^n) \quad (1.4)$$

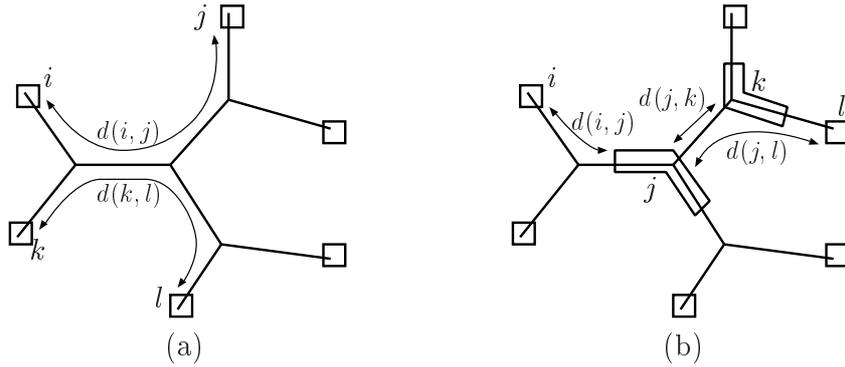


Figure 3: (a) tree metric and (b) distance between subtrees on a tree

where $c_{a,r}^f$ is the nonnegative number defined by $\sup\{t \geq 0 \mid f - t|\langle a, x \rangle - r| \text{ is convex}\}$ and f' is a polyhedral convex function satisfying $c_{a,r}^{f'} \in \{0, +\infty\}$ for all $(a, r) \in \mathbf{R}^n \times \mathbf{R}$. We call function $|\langle a, x \rangle - r|$ a *split function* corresponding to a split metric. Indeed, the terms of right hand side of Figure 2 can be regarded as the negative of split functions. We show in Section 2 that this decomposition (1.4) is uniquely determined (Theorem 2.2). We call this the *polyhedral split decomposition*.

As with the above example, this decomposition yields the decomposition of a function defined on a finite point set X . In Section 3, we study the split decomposition for a function defined on X (*the split decomposition of discrete functions*).

In Section 4, we regard a metric d as a function on Λ as Figure 1 and apply the results of Sections 2 and 3. We then obtain the following:

- By discretizing the polyhedral split decomposition (1.4), Bandelt-Dress' split decomposition (1.1) can be derived; see Proposition 4.5, Theorems 4.6, 4.7.
- The split decomposition of metrics can be naturally extended for distance functions, which may possibly violate the triangle inequality, using *partial split distances*; see Theorem 4.7, (4.13), and (4.27). This *extended split decomposition* exploits more detailed combinatorial structure of a given distance matrix than Bandelt-Dress' split decomposition (1.1) in the case that d violates the triangle inequality; see Remark 4.8. For example, a distance on subtrees on a tree defined by the shortest path length between subtrees (Figure 3), studied by [15], is a totally split-decomposable distance in our sense, which can be regarded as an extension of the fact that a tree metric is a totally split-decomposable metric in Bandelt-Dress' sense; $d' = 0$ in (1.1). So such distances can be recognized by our extended split decomposition; see Remark 4.18.
- We show *the split-hyperplane correspondence*, where each split can be associated with a certain hyperplane dividing simplex and weak compatibility of splits can be translated into a certain geometric property of the corresponding hyperplane arrangement and the point configuration Λ ; see Theorem 4.14 and Proposition 4.16.

Closing this introduction, we would like to mention some background of the present work. Applications of convex analysis to combinatorial and discrete structures are given in the theory of submodular functions by Lovász [19], Frank [11], and Fujishige [12] (see also [13]). Recently Murota [20] developed *discrete convex analysis*, a convex analysis for functions defined on integer lattice points. Hirai and Murota [16] discuss the relationship between tree metrics and discrete convexity, where they observe that d is a tree metric

if and only if the function on Λ associated with d , as Figure 1, is an *M-concave function*, which is one of fundamental class of discrete convex functions. The present paper is strongly influenced by these works.

2 Split decomposition of polyhedral convex functions

In this section, we derive the polyhedral split decomposition (1.4), which is the basis for subsequent developments in this paper.

2.1 Basic notation

Here we introduce some basic notation about convexity and polyhedra; see [22], [21], [27]. Let \mathbf{R} , \mathbf{R}_+ , \mathbf{R}_{++} be the sets of real numbers, nonnegative real numbers, and positive real numbers, respectively. Let \mathbf{R}^n be the n dimensional Euclidean space with the standard inner product $\langle \cdot, \cdot \rangle$. The norm $\|x\|$ of $x \in \mathbf{R}^n$ is defined by $\sqrt{\langle x, x \rangle}$. For $x, y \in \mathbf{R}^n$, let $[x, y]$ denote the closed line segment between x and y . For two subset $F, G \subseteq \mathbf{R}^n$, $F + G$ means the *Minkowski sum* of F and G defined by $\{x + y \mid x \in F, y \in G\}$. A subset $S \subseteq \mathbf{R}^n$ is said to be *convex* if $[x, y] \subseteq S$ for each $x, y \in S$, and is said to be (*convex*) *cone* if $\alpha x + \beta y \in S$ for $x, y \in S$ and $\alpha, \beta \in \mathbf{R}_+$. For a subset $X \subseteq \mathbf{R}^n$, we denote by $\text{conv } X$, $\text{cone } X$, and $\text{aff } X$ the convex hull (the smallest convex set containing X), the conical hull (the smallest convex cone containing X), and the affine hull (the smallest affine space containing X), respectively, i.e.,

$$\text{conv } X = \left\{ \sum_{y \in Y} \lambda_y y \mid Y \subseteq X : \text{finite set}, \lambda \in \mathbf{R}_+^Y, \sum_{y \in Y} \lambda_y = 1 \right\}, \quad (2.1)$$

$$\text{cone } X = \left\{ \sum_{y \in Y} \lambda_y y \mid Y \subseteq X : \text{finite set}, \lambda \in \mathbf{R}_+^Y \right\}, \quad (2.2)$$

$$\text{aff } X = \left\{ \sum_{y \in Y} \lambda_y y \mid Y \subseteq X : \text{finite set}, \lambda \in \mathbf{R}^Y, \sum_{y \in Y} \lambda_y = 0 \right\}. \quad (2.3)$$

$\text{int } X$ and $\text{ri } X$ mean the set of interior points and relative interior points of X , respectively, i.e., $\text{int } X = \{x \in X \mid \exists \epsilon > 0, \{x\} + B_\epsilon \subseteq X\}$ and $\text{ri } X = \{x \in X \mid \exists \epsilon > 0, \{x\} + \text{aff } X \cap B_\epsilon \subseteq X\}$, where $B_\epsilon = \{x \in \mathbf{R}^n \mid \|x\| \leq \epsilon\}$. In particular, for $(a, r) \in \mathbf{R}^n \times \mathbf{R}$, we define a hyperplane $H_{a,r}$ by $\{x \in \mathbf{R}^n \mid \langle a, x \rangle = r\}$, closed half spaces $H_{a,r}^-$ and $H_{a,r}^+$ by $\{x \in \mathbf{R}^n \mid \langle a, x \rangle \leq r\}$ and $\{x \in \mathbf{R}^n \mid \langle a, x \rangle \geq r\}$, and open half spaces $H_{a,r}^{--}$ and $H_{a,r}^{++}$ by $\{x \in \mathbf{R}^n \mid \langle a, x \rangle < r\}$ and $\{x \in \mathbf{R}^n \mid \langle a, x \rangle > r\}$. A subset $P \subseteq \mathbf{R}^n$ is a *polyhedron* if P is the intersection of a finite number of closed half spaces. A function $f : \mathbf{R}^n \rightarrow \mathbf{R} \cup \{+\infty\}$ is said to be convex if it satisfies

$$\lambda f(x) + (1 - \lambda)f(y) \geq f(\lambda x + (1 - \lambda)y) \quad (2.4)$$

for all $x, y \in \mathbf{R}^n$, $\lambda \in [0, 1]$. $\text{dom } f = \{x \in \mathbf{R}^n \mid f(x) < +\infty\}$ is called the *effective domain* of f and $\text{epi } f = \{(x, r) \in \mathbf{R}^n \times \mathbf{R} \mid f(x) \leq r\}$ is called the *epigraph* of f . The *indicator function* $\delta_S : \mathbf{R}^n \rightarrow \mathbf{R} \cup \{+\infty\}$ for a subset $S \subseteq \mathbf{R}^n$ is defined by

$$\delta_S(x) = \begin{cases} 0 & \text{if } x \in S \\ +\infty & \text{otherwise} \end{cases} \quad (x \in \mathbf{R}^n). \quad (2.5)$$

A function f is called a *polyhedral convex function* if its epigraph $\text{epi } f$ is polyhedral, or equivalently, f is represented by $\max\{\langle a_i, \cdot \rangle + b_i \mid 1 \leq i \leq m\} + \delta_P$ for $\{(a_i, b_i)\}_{i=1}^m \subseteq \mathbf{R}^n \times \mathbf{R}$ and a polyhedron P .

A nonempty subset F of a polyhedron P is called a *face* of P if $F = P$ or F can be represented as $H \cap P$ for some hyperplane H with $P \subseteq H^+$.

A *polyhedral complex* \mathcal{C} is a finite collection of polyhedra such that (1) if $P \in \mathcal{C}$, all the faces of P are also in \mathcal{C} , and (2) the nonempty intersection $P \cap Q$ of two polyhedra

$P, Q \in \mathcal{C}$ is a face of P and Q . The underlying set of \mathcal{C} is the point set $|\mathcal{C}| = \bigcup_{P \in \mathcal{C}} P$. A *polyhedral subdivision* of a polyhedron P is a polyhedral complex \mathcal{C} with $|\mathcal{C}| = P$.

For a polyhedral convex function f , let $\mathcal{T}(f)$ be the polyhedral subdivision of $\text{dom } f$ induced by the projection of lower faces $\text{epi } f$, which is represented as

$$\mathcal{T}(f) = \{F \subseteq \mathbf{R}^n \mid F \text{ is the set of minimizers of } f(\cdot) - \langle p, \cdot \rangle \text{ for some } p \in \mathbf{R}^n\}. \quad (2.6)$$

Such subdivisions are said to be *regular*; see [27, Chapter 5]. Note that on each $F \in \mathcal{T}(f)$, f is an affine function.

For two polyhedral subdivisions \mathcal{C}_1 and \mathcal{C}_2 , the *common refinement* $\mathcal{C}_1 \wedge \mathcal{C}_2$ is defined by $\mathcal{C}_1 \wedge \mathcal{C}_2 = \{F \cap G \mid F \in \mathcal{C}_1, G \in \mathcal{C}_2, F \cap G \neq \emptyset\}$. Note that $\mathcal{C}_1 \wedge \mathcal{C}_2$ is a polyhedral subdivision of $|\mathcal{C}_1| \cap |\mathcal{C}_2|$. In particular, for a set of hyperplanes \mathcal{H} , we define the polyhedral subdivision $\mathcal{A}(\mathcal{H})$ of \mathbf{R}^n as

$$\mathcal{A}(\mathcal{H}) = \bigwedge_{H \in \mathcal{H}} \{H, H^+, H^-\}. \quad (2.7)$$

Namely, $\mathcal{A}(\mathcal{H})$ is the partition of \mathbf{R}^n by hyperplanes \mathcal{H} . We give a fundamental relationship between $\mathcal{T}(f + g)$ and $\mathcal{T}(f) \wedge \mathcal{T}(g)$.

Lemma 2.1. *Let f and g be polyhedral convex functions with $\text{dom } f \cap \text{dom } g \neq \emptyset$. Then we have*

$$\mathcal{T}(f + g) = \mathcal{T}(f) \wedge \mathcal{T}(g). \quad (2.8)$$

Proof. Both sides of (2.8) are polyhedral subdivisions of $\text{dom } f \cap \text{dom } g$. Hence it is sufficient to show the inclusion (\supseteq). Let $F \in \mathcal{T}(f)$ and $G \in \mathcal{T}(g)$ be the sets of minimizers of $f(\cdot) - \langle p, \cdot \rangle$ and $g(\cdot) - \langle q, \cdot \rangle$ for some $p, q \in \mathbf{R}^n$, respectively. If $F \cap G \neq \emptyset$, then $F \cap G$ is the set of minimizers of $(f + g)(\cdot) - \langle p + q, \cdot \rangle$. Hence $F \cap G \in \mathcal{T}(f + g)$. \square

2.2 Polyhedral split decomposition

Here we derive the split decomposition of polyhedral convex functions, which is our guiding principle in deriving the split decomposition of metrics and related concepts in terms of convex analysis and polyhedral geometry. Furthermore our derivation leads to an algorithm for the discrete version of polyhedral split decomposition in the next section.

First we define a *split function*, which is a correspondence of a split metric. For a hyperplane $H = H_{a,b}$ with $\|a\| = 1$, the split function $l_H : \mathbf{R}^n \rightarrow \mathbf{R}$ associated with H is defined as

$$l_H(x) = |\langle a, x \rangle - b|/2 \quad (x \in \mathbf{R}^n). \quad (2.9)$$

For a polyhedral convex function f and a hyperplane H , we define the nonnegative number $c_H(f) \in \mathbf{R}_+ \cup \{+\infty\}$ as

$$c_H(f) = \sup\{t \in \mathbf{R}_+ \mid f - tl_H \text{ is convex}\} \quad (2.10)$$

and define the set of hyperplane $\mathcal{H}(f)$ as

$$\mathcal{H}(f) = \{H : \text{hyperplane} \mid 0 < c_H(f) < +\infty\}. \quad (2.11)$$

Note that $c_H(f) = +\infty$ if and only if $\text{dom } f \subseteq H^+$ or $\text{dom } f \subseteq H^-$. Main aim here is to prove the following.

Theorem 2.2. A polyhedral convex function $f : \mathbf{R}^n \rightarrow \mathbf{R} \cup \{+\infty\}$ with $\dim \text{dom } f = n$ can be decomposed as

$$f = \sum_{H \in \mathcal{H}(f)} c_H(f) l_H + f', \quad (2.12)$$

where f' is a polyhedral convex function with $c_H(f') \in \{0, +\infty\}$ for each hyperplane H . Moreover, this representation is unique in the following sense. If f is represented as

$$f = \sum_{H \in \mathcal{H}} \alpha_H l_H + g \quad (2.13)$$

for some set of hyperplanes \mathcal{H} each of whose elements meets $\text{int dom } f$, positive numbers $\{\alpha_H\}_{H \in \mathcal{H}}$, and a polyhedral convex function g with $c_H(g) \in \{0, +\infty\}$ for any hyperplane H , then we have

$$\mathcal{H} = \mathcal{H}(f), \quad \alpha_H = c_H(f) \quad (H \in \mathcal{H}), \quad f' = g. \quad (2.14)$$

The proof needs three lemmas (Lemma 2.3, 2.4, 2.5). The first is a convexity condition of a function $f - tl_H$. The second is a condition of $0 < c_H(f) < +\infty$ for f and H . The third is a behavior of $c_{H'}(f - tl_H)$.

The proof uses the structure of the subdivision $\mathcal{T}(f)$ defined by (2.6). In particular, a full dimensional polyhedron of $\mathcal{T}(f)$ is called a *linearity domain* of f . Note that such members exist since $\text{dom } f$ is full dimensional. A function $g : \mathbf{R}^n \rightarrow \mathbf{R} \cup \{+\infty\}$ is called a $\mathcal{T}(f)$ -*piecewise-linear function* if $\text{dom } g = \text{dom } f$, g is continuous on $\text{dom } g$, and is an affine function on each linearity domain of f . We shall discuss the convexity condition for $\mathcal{T}(f)$ -piecewise-linear function g . On each linearity domain F , g can be uniquely represented as

$$g(x) = \langle p_F^g, x \rangle - b_F^g \quad (x \in F) \quad (2.15)$$

for some $(p_F^g, b_F^g) \in \mathbf{R}^n \times \mathbf{R}$. A pair of linearity domains F and G is said to be *adjacent* if $\dim F \cap G = n - 1$. In particular, we denote the hyperplane $\text{aff } F \cap G$ by $H_{F,G}$ and assume $F \subseteq H_{F,G}^+$ and $G \subseteq H_{F,G}^-$. For two adjacent linearity domains $F, G \in \mathcal{T}(f)$, we define a function $g_{F,G}$ as

$$g_{F,G}(x) = \begin{cases} \langle p_F^g, x \rangle - b_F^g & \text{if } x \in H_{F,G}^+ \\ \langle p_G^g, x \rangle - b_G^g & \text{if } x \in H_{F,G}^- \end{cases} \quad (x \in \mathbf{R}^n) \quad (2.16)$$

In particular, g coincides with $g_{F,G}$ on $F \cup G$. By continuity of g , $g_{F,G}$ can be also represented by

$$g_{F,G} = \begin{cases} \max\{\langle p_F^g, \cdot \rangle - b_F^g, \langle p_G^g, \cdot \rangle - b_G^g\} & \text{if } g_{F,G} \text{ is convex} \\ \min\{\langle p_F^g, \cdot \rangle - b_F^g, \langle p_G^g, \cdot \rangle - b_G^g\} & \text{if } g_{F,G} \text{ is concave} \end{cases} \quad (2.17)$$

$$= \begin{cases} \|p_F^g - p_G^g\| l_{H_{F,G}} + \text{affine function} & \text{if } g_{F,G} \text{ is convex} \\ -\|p_F^g - p_G^g\| l_{H_{F,G}} + \text{affine function} & \text{if } g_{F,G} \text{ is concave} \end{cases} \quad (2.18)$$

Note that if $p_F^g \neq p_G^g$, then $H_{F,G} = H_{p_F^g - p_G^g, b_F^g - b_G^g}$.

Lemma 2.3. A $\mathcal{T}(f)$ -piecewise-linear function g is convex if and only if $g_{F,G}$ is convex for each pair of adjacent linearity domains F and G of f .

Proof. Only-if part is obvious by $g = g_{F,G}$ on $F \cup G$. We show if part. For sufficiently generic points $x, y \in \text{dom } g$, the line segment $[x, y]$ meets only n or $(n - 1)$ -dimensional members of $\mathcal{T}(f)$. By condition, g is convex on $[x, y]$, i.e., g satisfies (2.4) for x, y and $0 \leq \lambda \leq 1$. Hence, by continuity of g , g satisfies (2.4) for any $x, y \in \text{dom } f$. \square

Lemma 2.4. For a hyperplane $H = H_{a,b}$ for $\|a\| = 1$, we have $0 < c_H(f) < +\infty$ if and only if $H \cap \text{int dom } f \neq \emptyset$ and $F \subseteq H^+$ or $F \subseteq H^-$ for each $F \in \mathcal{T}(f)$.

Proof. Only-if part is easily observed. Indeed, if there exists $F \in \mathcal{T}(f)$ such that both $F \cap H^{++}$ and $F \cap H^{--}$ are nonempty, then $f - tl_H$ is not convex on F for any $t > 0$. We show if part. By the condition, $f - tl_H$ is a $\mathcal{T}(f)$ -piecewise-linear function. By Lemma 2.3, $f - tl_H$ is convex if and only if $(f - tl_H)_{F,G}$ is convex for each pair of adjacent linearity domains F, G of f . If $F, G \subseteq H^+$ or $F, G \subseteq H^-$, then we have

$$(f - tl_H)_{F,G} = f_{F,G} + \text{affine function} \quad (2.19)$$

and $(f - tl_H)_{F,G}$ is convex for $t \geq 0$. If $F \subseteq H^+$ and $G \subseteq H^-$, then we have $\text{aff } F \cap G = H_{F,G} = H_{a,b}$ and $p_F^f - p_G^f = \pm \|p_F^f - p_G^f\|a$. Hence we have

$$(f - tl_H)_{F,G} = (\|p_F^f - p_G^f\| - t)l_H + \text{affine function.} \quad (2.20)$$

Hence we have

$$c_H(f) = \min\{\|p_F^f - p_G^f\| \mid F, G \text{ are linearity domains with } \text{aff } F \cap G = H\} \quad (2.21)$$

By definition of linearity domain, we have $p_F^f \neq p_G^f$ and $\|p_F^f - p_G^f\| > 0$. By polyhedrality of f , pairs of adjacent linearity domains are finite. Hence we have $0 < c_H < +\infty$. \square

Lemma 2.5. For $H, H' \in \mathcal{H}(f)$ and $t \in [0, c_H(f)]$, we have

$$c_{H'}(f - tl_H) = \begin{cases} c_H(f) - t & \text{if } H' = H, \\ c_{H'}(f) & \text{otherwise.} \end{cases} \quad (2.22)$$

Proof. For distinct $H, H' \in \mathcal{H}(f)$, two subsets

$$\{\{F, G\} : \text{a pair of linearity domains of } f \text{ with } \text{aff } F \cap G = H\}, \quad (2.23)$$

$$\{\{F, G\} : \text{a pair of linearity domains of } f \text{ with } \text{aff } F \cap G = H'\} \quad (2.24)$$

are disjoint. Hence, (2.22) follows from (2.19), (2.20), and (2.21) in the proof of Lemma 2.4. \square

The proof of Theorem 2.2. First note that $\mathcal{H}(f)$ is necessarily finite since H is represented by $\text{aff } F \cap G$ for some adjacent linearity domains F and G of f and the set of linearity domains are finite by polyhedrality of f . By Lemma 2.5, $f' = f - \sum_{H \in \mathcal{H}(f)} c_H(f)l_H$ is convex and $c_H(f') \in \{0, +\infty\}$ for any hyperplane H . Hence we obtain the first part. We show the uniqueness part. Suppose that f can be decomposed as the form (2.13). Then we have $\mathcal{H} \subseteq \mathcal{H}(f)$ and $\alpha_H \leq c_H(f)$ for $H \in \mathcal{H}$. By Lemma 2.5, we have

$$0 = c_{H'}(g) = c_{H'}(f - \sum_{H \in \mathcal{H}} \alpha_H l_H) = \begin{cases} c_H(f) - \alpha_H & \text{if } H' = H \text{ for some } H \in \mathcal{H} \\ c_{H'}(f) & \text{otherwise} \end{cases} \quad (2.25)$$

for $H' \in \mathcal{H}(f)$. This implies (2.14). \square

Remark 2.6. We observe that $\mathcal{T}(\alpha l_H) = \{H, H^+, H^-\}$ for $\alpha \in \mathbf{R}_{++}$. Hence, by Lemma 2.1, corresponding to the decomposition (2.12), $\mathcal{T}(f)$ is decomposed as

$$\mathcal{T}(f) = \mathcal{A}(\mathcal{H}(f)) \wedge \mathcal{T}(f'). \quad (2.26)$$

Then, by Lemma 2.4, there is no hyperplane H meeting $\text{int dom } f$ such that $F \subseteq H^+$ and $F \subseteq H^-$ for each $F \in \mathcal{T}(f')$. This gives a geometric interpretation of the split-primeness of metrics in Section 4.

Closing this subsection, we give an alternative expression of $c_H(f)$, which will be used in Section 3.

Lemma 2.7. *For a polyhedral convex function $f : \mathbf{R}^n \rightarrow \mathbf{R} \cup \{+\infty\}$ with $\dim \text{dom } f = n$ and a hyperplane H , we have*

$$c_H(f) = \frac{1}{2} \inf \left\{ \frac{f(x) - f(w)}{l_H(x)} + \frac{f(y) - f(w)}{l_H(y)} \mid \begin{array}{l} x \in \text{dom } f \cap H^{++} \\ y \in \text{dom } f \cap H^{--} \\ \{w\} = [x, y] \cap H \end{array} \right\}. \quad (2.27)$$

In particular, if $c_H(f) = 0$, then there exists some linearity domain F of f such that both $F \cap H^{++}$ and $F \cap H^{--}$ are nonempty and any $x \in F \cap H^{++}$ and $y \in F \cap H^{--}$ attain the infimum of RHS of (2.27). If $0 < c_H(f) < +\infty$, then there exists an adjacent pair of linearity domain $F \subseteq H^+$, $G \subseteq H^-$ of f such that $\text{aff}(F \cap G) = H$ and any $x \in F \cap H^{++}$ and $y \in G \cap H^{--}$ attain the infimum of RHS of (2.27).

Proof. First we show $c_H(f) \leq \text{RHS of (2.27)}$. Indeed, since $f - c_H(f)l_H$ is convex, $f - c_H(f)l_H$ satisfies (2.4) for $x \in H^{++}$, $y \in H^{--}$, and $\lambda = l_H(y)/(l_H(x) + l_H(y))$ as

$$\lambda(f - c_H(f)l_H)(x) + (1 - \lambda)(f - c_H(f)l_H)(y) \leq (f - c_H(f))(\lambda x + (1 - \lambda)y). \quad (2.28)$$

Put $w = \lambda x + (1 - \lambda)y$. Then we have $\{w\} = H \cap [x, y]$ and $l_H(w) = 0$. Hence we obtain

$$c_H(f) \leq \frac{1}{2} \left\{ \frac{f(x) - f(w)}{l_H(x)} + \frac{f(y) - f(w)}{l_H(y)} \right\}. \quad (2.29)$$

Therefore, it suffices to show the latter part. The statement for the case $c_H(f) = 0$ follows from Lemma 2.4. In the case $0 < c_H(f) < +\infty$, by (2.18) and (2.21), there exists a pair of adjacent linearity domains $F \subseteq H^+$, $G \subseteq H^-$ of f such that

$$f_{F,G} = c_H(f)l_H + \text{affine function}. \quad (2.30)$$

Furthermore, $F \cup G$ is a linearity domain of $f - c_H(f)l_H$ and therefore is convex. Then, for any $x \in F \cap H^{++}$, $y \in G \cap H^{--}$, we have $\{w\} = [x, y] \cap H \in F \cup G$ and

$$\begin{aligned} \text{RHS of (2.29)} &= \frac{1}{2} \left\{ \frac{f_{F,G}(x) - f_{F,G}(w)}{l_H(x)} + \frac{f_{F,G}(y) - f_{F,G}(w)}{l_H(y)} \right\} \\ &= c_H(f). \end{aligned} \quad (2.31)$$

□

2.3 Dual view of Theorem 2.2

In this subsection, we discuss the dual view of Theorem 2.2. In particular, we discuss relationship between Theorem 2.2 and the maximum zonotopic summand of a polyhedron. For this, we focus on the case that f is positively homogeneous, or equivalently, is the support function of some polyhedron. Recall the relationship between positively homogeneous convex functions and the support functions of a polyhedron; see [22, Section 1.7], [21, Section 19].

For a polyhedron P , the *support function* f_P is defined as

$$f_P(x) = \sup\{\langle x, p \rangle \mid p \in P\} \quad (x \in \mathbf{R}^n). \quad (2.32)$$

Then f_P is polyhedral convex and *positively homogeneous*, i.e., f satisfies

$$f(\lambda x) = \lambda f(x) \quad \lambda \in \mathbf{R}_+, x \in \mathbf{R}^n. \quad (2.33)$$

Conversely, if f is positively homogeneous polyhedral convex, then f is the support function of the polyhedron

$$P_f = \{p \in \mathbf{R}^n \mid \langle p, x \rangle \leq f(x) \ (x \in \mathbf{R}^n)\}. \quad (2.34)$$

This correspondence is one to one and satisfies

$$P_{f+g} = P_f + P_g, \quad f_{P+Q} = f_P + f_Q \quad (2.35)$$

for two positively homogeneous polyhedral convex functions f and g and two polyhedra P and Q . Note that the split function l_H for a linear hyperplane $H = H_{a,0}$ with $\|a\| = 1$ is the support function of the line segment $(1/2)[-a, a]$.

In this case, $\mathcal{T}(f_P)$ coincides with the *normal fan* of P , which consists of the *normal cones* of P_f . Recall the definition of the normal cone; see [27, Chapter 7], [22, Section 2.4]. For a face F of P_f , the normal cone N_F of F is defined as

$$N_F = \{x \in \mathbf{R}^n \mid \{q \in P \mid \langle x, q \rangle = \sup\{\langle x, p \rangle \mid p \in P\}\} \subseteq F\} \quad (2.36)$$

and hence we have $\mathcal{T}(f_P) = \{N_F \mid F \text{ is a face of } P\}$. Note that $\text{dom } f_P$ is given by $\{x \in \mathbf{R}^n \mid \langle p, x \rangle \leq 0 \ (p \in \text{rec } P)\}$, where $\text{rec } P$ is the *recession cone* of P defined by $\{v \in \mathbf{R}^n \mid P + [0, v] = P\}$.

In particular, faces of P_f and $\mathcal{T}(f)$ is in polar relation as follows.

(N1) $F \subseteq G$ if and only if $N_F \supseteq N_G$, and

(N2) the linear subspace parallel to $\text{aff } F$ is orthogonal to $\text{aff } N_F$,

where F and G are faces of P_f . In addition, we see the following relation.

(N3) F is bounded if and only if $N_F \cap \text{int dom } f \neq \emptyset$.

Applying Theorem 2.2 to the support function of a polyhedron P , we have the following, which can be understood as an unbounded refinement of theorem of Bolker [4, p.341] that every polytope (bounded polyhedron) has the maximum zonotopic summand, where a polyhedron Q is called a *Minkowski summand* of a polyhedron P if there exists a polyhedron P' such that $P = Q + P'$, and P is said to be *pointed* if P has a vertex.

Corollary 2.8. *A pointed polyhedron P is uniquely decomposed as*

$$P = \sum_{v \in \mathcal{V}} [-v/2, v/2] + P' \quad (2.37)$$

for a finite set of vectors $\mathcal{V} \subseteq \mathbf{R}^n \setminus \pm \text{rec } P$ and a polyhedron P' which does not contain $[-u/2, u/2]$ as Minkowski summand for each $u \in \mathbf{R}^n \setminus \pm \text{rec } P$

Proof. Let f_P be the support function of P . Note that $\text{dom } f_P$ is full dimensional since P is pointed. Since $\mathcal{T}(f)$ consists of polyhedral cones, $\mathcal{H}(f)$ consists of linear hyperplanes. From $\text{dom } f_P = \{x \in \mathbf{R}^n \mid \langle p, x \rangle \leq 0 \ (p \in \text{rec } P)\}$, we have $\text{dom } f_P \subseteq H_{a,0}^+$ or $\text{dom } f_P \subseteq H_{a,0}^-$ if and only if $a \in \text{rec } P$ or $-a \in \text{rec } P$. Applying Theorem 2.2 to f_P and (2.35), we obtain the desired result. \square

A finite (Minkowski) sum of line segments is called a *zonotope*. The zonotope $Z(P) = \sum_{v \in \mathcal{V}} [-v/2, v/2]$ of (2.37) is called the *maximum zonotopic summand* of P . As will be seen in Section 4, the split decomposition of metrics can be understood as the extraction of the maximum zonotopic summand from some polyhedron associated with a given metric; see Remark 4.10.

The structure of the union of the bounded faces of P_f expresses the nonlinearity of f over $\text{dom } f$. For example, f is linear on $\text{dom } f$ if and only if the union of the bounded faces of P_f is a single point. As will be shown in Section 4, the following proposition gives an interpretation and another proof to the result of Dress [7] that a metric is a tree metric if and only if its tight span is a tree.

Proposition 2.9. *For a positively homogeneous polyhedral convex function $f : \mathbf{R}^n \rightarrow \mathbf{R} \cup \{+\infty\}$ with $\dim \text{dom } f = n$, the following conditions are equivalent.*

- (1) *The dimension of the union of bounded faces of P_f is (less than or) equal to 1.*
- (2) *The union of bounded faces of P_f is a tree.*
- (3) *If f is decomposed as (2.12), then f' is linear over $\text{dom } f$ and $H_1 \cap H_2 \cap \text{int } \text{dom } f = \emptyset$ for each pair $H_1, H_2 \in \mathcal{H}(f)$.*

Proof. The equivalence (1) \Leftrightarrow (2) follows from the general fact that the union of bounded faces of a polyhedron is contractible, i.e., it is homotopy equivalent to a point (see [15, Appendix] for a proof). We show (3) \Rightarrow (1). By condition, each member of $\mathcal{T}(f)$ meeting $\text{int } \text{dom } f$ has the dimension n or $n - 1$. By (N2), we have (1). Finally we show (2) \Rightarrow (3). Let E be a 1-dimensional bounded face of P_f . Then F is a line segment. By (N2), N_E is an $(n - 1)$ -dimensional member of $\mathcal{T}(f)$. We claim that $\text{aff } N_E \cap \text{dom } f_P = N_E$. Indeed, this follows from that fact every proper faces of N_E are contained in the boundary of $\text{dom } f$ by (N2). By Lemma 2.4, we conclude that (linear) hyperplanes $\{\text{aff } N_E \mid E \text{ is a bounded edge of } P\} = \mathcal{H}(f)$ and do not intersect in $\text{int } \text{dom } f$ each other. Moreover, N_E can be uniquely represented by $F \cap G$ for two adjacent linearity domains F and G . Then $F \cup G$ is a linearity domain of $f - c_H l_H$. Repeat this process to $f - c_H(f) l_H$, then $\text{dom } f$ is a linearity domain of $f - \sum_{H \in \mathcal{H}(f)} c_H(f) l_H$. This implies that $f - \sum_{H \in \mathcal{H}(f)} c_H(f) l_H$ is linear over $\text{dom } f$. \square

3 Split decomposition of discrete functions

In this section, we describe a discrete version of the polyhedral split decomposition, which is a certain kind of decomposition of a function f defined on a finite set X of points of \mathbf{R}^n . Basic idea is to associate f to some polyhedral convex function \bar{f} as Figure 1 in the introduction and apply Theorem 2.2.

3.1 Discrete functions and convex-extension

Let X be a finite point set of \mathbf{R}^n . We assume that X is contained by some affine hyperplane K not containing origin of \mathbf{R}^n . If necessary, we consider $\tilde{X} = \{(y, 1) \mid y \in X\} \subseteq \mathbf{R}^{n+1}$ instead of X ; in this case $K = \{(x, 1) \mid x \in \mathbf{R}^n\} \subseteq \mathbf{R}^{n+1}$. For a function $f : X \rightarrow \mathbf{R}$, we consider the following polyhedral convex function

$$\bar{f}(x) = \sup \left\{ \sum_{y \in X} \lambda_y f(y) \mid \sum_{y \in X} \lambda_y \chi_y = x, \lambda_y \geq 0 (y \in X) \right\}, \quad (3.1)$$

which is called the *homogeneous convex extension* of f . Then, by linear programming duality, \bar{f} is the support function of the polyhedron

$$Q(f) = \{p \in \mathbf{R}^n \mid \langle p, y \rangle \leq f(y) (y \in X)\} \quad (3.2)$$

and, therefore, positively homogeneous. $\text{dom } \bar{f}$ is given by $\text{cone } X$. Then $\mathcal{T}(\bar{f})$ is the normal fan of $Q(f)$. Let $\mathcal{T}^X(\bar{f})$ denote the subdivision of $\text{conv } X$ which is defined by

$\{F \cap K \mid F \in \mathcal{T}(\bar{f})\}$. Note that the restriction of \bar{f} to K is the function of the lower envelope of the convex hull of $\{(y, f(y)) \mid y \in X\} \subseteq K \times \mathbf{R}$ and $\mathcal{T}^X(\bar{f})$ is the projection of lower faces of $\text{conv}\{(y, f(y)) \mid y \in X\}$; recall Figure 1 in the introduction. For a function $g : \mathbf{R}^n \rightarrow \mathbf{R} \cup \{+\infty\}$, we denote the restriction of g to X by g^X . f is said to be *convex-extensible* if it satisfies $\bar{f}^X = f$. The set of convex-extensible functions is recognized as a fundamental class of the *discrete convex functions* (see [20]).

We give some fundamental properties of discrete functions and homogeneous convex-extensions, which will be used in the proof of Theorem 3.4 and Proposition 3.6 in the next subsection.

Lemma 3.1. *Let $f, g : X \rightarrow \mathbf{R}$ be a function.*

- (1) $\overline{cf + \langle q, \cdot \rangle} = c\bar{f} + \langle q, \cdot \rangle + \delta_{\text{cone } X}$ ($c \in \mathbf{R}_+, q \in \mathbf{R}^n$), where $\delta_{\text{cone } X}$ is the indicator function of cone X ; see (2.5).
- (2) $F \in \mathcal{T}(\bar{f})$ is represented as $\text{cone}\{y \mid y \in X, \langle q, y \rangle = f(y)\}$ for some $q \in Q(f)$. Furthermore $\bar{f}(x) = \overline{f^{F \cap X}}(x)$ for $x \in F$. In particular, $\bar{f}(y) = f(y)$ if $y \in X$ is a vertex of $\mathcal{T}^X(\bar{f})$.

Proof. (1). By (3.2), we have

$$\overline{cf + \langle q, \cdot \rangle}(x) = \sup\{\langle x, p \rangle \mid p \in \mathbf{R}^n, \langle p, y \rangle \leq cf(y) + \langle q, y \rangle \ (y \in X)\}. \quad (3.3)$$

If $c > 0$, substitute $p' = (p - q)/c$ to RHS of (3.3). Then we have the desired result. If $c = 0$, then it is immediate from the definition (3.1).

(2). The first statement follows from the fact that a normal cone of $Q(f)$ can be represented as the conical combination of the normal vectors of facets of $Q(f)$ containing some $q \in Q(f)$. Next we show $\bar{f}(x) = \overline{f^{F \cap X}}(x)$ for $x \in F$. Indeed, the coefficient λ of conical combination $x = \sum_{y \in X, \langle q, y \rangle = f(y)} \lambda_y y$ and q satisfies the feasibility and the complementary slackness condition for the dual pairs of linear programs corresponding to $\bar{f}(x)$ and $\overline{f^{F \cap X}}(x)$. Hence we obtain $\bar{f}(x) = \overline{f^{F \cap X}}(x) = \sum_{y \in X, \langle q, y \rangle = f(y)} \lambda_y f(y)$; see [23] for linear programming. \square

3.2 Split decomposition of discrete functions

Here, we present the discrete version of Theorem 2.2. As shown in the next section, this turns out to coincide with (an extension of) Bandelt-Dress' split decomposition in the case of $X = \Lambda$. Furthermore some interesting properties of weakly compatible splits are also generalized. Arguments in this subsection have own significance in the theory of regular subdivision; see Remark 3.8. Note that since $\mathcal{T}(\bar{f})$ for $f : X \rightarrow \mathbf{R}$ is the normal fan of $Q(f)$, each hyperplane $H \in \mathcal{T}(\bar{f})$ is linear, i.e., $H = H_{a,0}$ for some $a \in \mathbf{R}^n$. We assume that $\text{aff } X = K$, which guarantees the full dimensionality of $\text{dom } \bar{f} = \text{cone } X$ for $f \in \mathbf{R}^X$. The discrete version of Theorem 2.2 is as follows.

Theorem 3.2. *A discrete function $f : X \rightarrow \mathbf{R}$ can be decomposed as*

$$f = \sum_{H \in \mathcal{H}(\bar{f})} c_H(\bar{f}) l_H^X + \gamma, \quad (3.4)$$

where $\gamma : X \rightarrow \mathbf{R}$ satisfies $c_{H'}(\bar{\gamma}) \in \{0, +\infty\}$ for any linear hyperplane H' . Moreover we have

$$\bar{f} = \sum_{H \in \mathcal{H}(\bar{f})} c_H(\bar{f}) l_H + \bar{\gamma}, \quad (3.5)$$

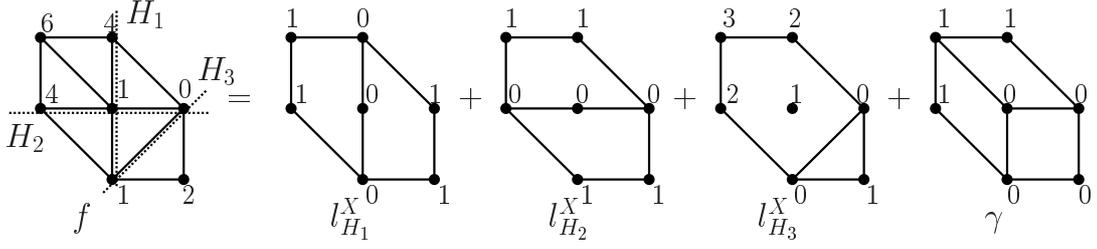


Figure 4: the split decomposition of f

or equivalently,

$$Q(f) = \sum_{H \in \mathcal{H}(\bar{f})} c_H(\bar{f})[-n_H/2, n_H/2] + Q(\gamma), \quad (3.6)$$

where n_H is the unit normal vector of hyperplane H . If, in addition, f is convex-extensible, then γ is also convex-extensible.

Proof. We can apply Theorem 2.2 to \bar{f} by $\dim \bar{f} = \dim \text{cone } X = n$. It suffices to show $\bar{f} = t l_H + (f - t l_H^X)$, or equivalently,

$$Q(f) = t[n_H/2, -n_H/2] + Q(f - t l_H^X) \quad (3.7)$$

for $H \in \mathcal{H}(\bar{f})$, and $t \in [0, c_H(\bar{f})]$. We show (3.7). The inclusion (\supseteq) follows from

$$\langle p + s n_H/2, y \rangle \leq f(y) - t |\langle n_H, y \rangle|/2 + s \langle n_H, y \rangle/2 \leq f(y) \quad (y \in X) \quad (3.8)$$

for $s \in [-t, t]$ and $p \in Q(f - t l_H^X)$. We next show (\subseteq). By $\bar{f} = t l_H + (\bar{f} - t l_H)$ and (2.35), we have

$$Q(f) = P_{\bar{f}} = t[n_H/2, -n_H/2] + P_{\bar{f} - t l_H}. \quad (3.9)$$

By the definition of $Q(\cdot)$ (3.2), (2.34), and $\bar{f}^X \leq f$, we have

$$P_{\bar{f} - t l_H} \subseteq Q((\bar{f} - t l_H)^X) \subseteq Q(f - t l_H^X). \quad (3.10)$$

Hence, we obtain (\subseteq). If f is convex-extensible, then $\bar{f}^X = f$. By restricting both side of (3.10) to X , we have $\gamma = \bar{\gamma}^X$. Hence γ is convex-extensible. \square

Figure 4 is an example of the split decomposition (3.4) of a function $f : X \rightarrow \mathbf{R}$ for $X \subseteq \mathbf{R}^3$ in terms of subdivisions induced by each terms in (3.4), where this figure illustrates the restriction to 2-dimensional space $K = \text{aff } X$, each black point denotes a point of X , and the number near each black point denotes corresponding value of f . The property (3.5)(or (3.6)) means that (3.4) decomposes f to the sum of functions each of which induces the subdivision coarser than $\mathcal{T}^X(\bar{f})$. This property corresponds to the coherence property of additive decomposition of metrics; see Section 4.

From Figure 4, we observe that possible hyperplanes appearing $\mathcal{H}(\cdot)$ are limited by the point set X . Motivated by this observation, we define a certain geometric condition of hyperplane arrangements, which is determined solely by X , as follows. A set of linear hyperplane \mathcal{H} is X -admissible if \mathcal{H} satisfies

(A0) $H \cap \text{ri conv } X \neq \emptyset$ for each $H \in \mathcal{H}$, and

(A1) $\text{conv}(F \cap X) = F \cap \text{conv } X$ for each $F \in \mathcal{A}(\mathcal{H})$.

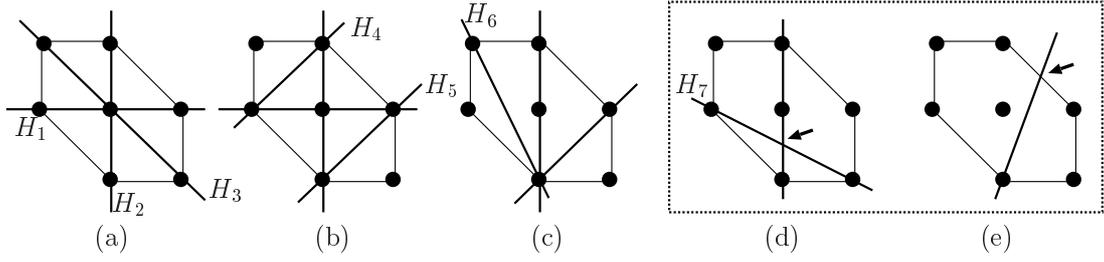


Figure 5: X -admissibility of hyperplanes

In particular, the condition (A1) can be rephrased as

$$\{\text{vertices of } (\mathcal{A}(\mathcal{H}) \wedge \{\text{faces of } \text{conv } X\})\} \subseteq X. \quad (3.11)$$

Namely, when \mathcal{H} cuts $\text{conv } X$, newly appeared vertices are contained in X .

Lemma 3.3. *For $f : X \rightarrow \mathbf{R}$, the set of hyperplanes $\mathcal{H}(\bar{f})$ is X -admissible.*

Proof. (A0) is clearly satisfied. We show (A1). The inclusion (\subseteq) is obvious. We show (\supseteq). By (2.26) and Lemma 3.1 (2), we have

$$\begin{aligned} F \cap \text{conv } X &= \bigcup \{G \in \mathcal{T}^X(\bar{f}) \mid G \subseteq F\} \\ &= \bigcup \{\text{conv}(G \cap X) \mid G \in \mathcal{T}^X(\bar{f}), G \subseteq F\} \\ &\subseteq \text{conv}(X \cap \bigcup \{G \in \mathcal{T}^X(\bar{f}) \mid G \subseteq F\}) = \text{conv}(X \cap F). \end{aligned} \quad (3.12)$$

□

Figure 5 illustrates examples of X -admissible and non X -admissible sets of hyperplanes, where black points are X , and hyperplanes of (a), (b), and (c) are X -admissible, and (d) and (e) are not since there are vertices, indicated by the arrow, not contained by X . Note that if \mathcal{H} is X -admissible, then any subset of \mathcal{H} is also X -admissible. So we define the set of linear hyperplanes \mathcal{H}_X as

$$\mathcal{H}_X = \{H : \text{linear hyperplane} \mid \{H\} \text{ is } X\text{-admissible}\}. \quad (3.13)$$

By (A0), (A1), and (3.11), we see that $H \in \mathcal{H}_X$ if and only if H satisfies

- (a0) there exists $x, y \in X$ such that $x \in H^{++}$ and $y \in H^{--}$, and
- (a1) for each edge $[x, y]$ of $\text{conv } X$ with $x \in H^{++}$, $y \in H^{--}$, we have $[x, y] \cap H \in X$.

Note that \mathcal{H}_X is a finite set since $H \in \mathcal{H}_X$ is represented as the linear space spanned by some $Y \subseteq X$. An X -admissible set of hyperplanes is a subset of \mathcal{H}_X . The next theorem implies that the discrete split decomposition (3.4) can be carried out without explicit construction of convex-extensions; $c_H(\bar{f})$ can be calculated discretely.

Theorem 3.4. *For a function $f : X \rightarrow \mathbf{R}$ and a hyperplane $H \in \mathcal{H}_X$, let $\tilde{c}_H(f)$ be defined by*

$$\frac{1}{2} \min \left\{ \frac{f(x) - \overline{f^{X \cap H}}(w)}{l_H(x)} + \frac{f(y) - \overline{f^{X \cap H}}(w)}{l_H(y)} \mid \begin{array}{l} x \in X \cap H^{++} \\ y \in X \cap H^{--} \\ \{w\} = H \cap [x, y] \end{array} \right\}. \quad (3.14)$$

Then we have

$$c_H(\bar{f}) = \max(0, \tilde{c}_H(f)). \quad (3.15)$$

Proof. By Lemma 2.7 and Lemma 3.1 (2), $c_H(\bar{f})$ can be represented as

$$c_H(\bar{f}) = \frac{1}{2} \min \left\{ \frac{f(y) - \bar{f}(w)}{l_H(y)} + \frac{f(z) - \bar{f}(w)}{l_H(z)} \mid \begin{array}{l} y \in X \cap H^{++} \\ z \in X \cap H^{--} \\ \{w\} = H \cap [x, y] \end{array} \right\}. \quad (3.16)$$

Since $\overline{f^{X \cap H}}(w) \geq \bar{f}(w)$ holds for $w \in H \cap \text{cone } X$, we have $c_H(\bar{f}) \geq \tilde{c}_H(f)$. If $c_H(\bar{f}) > 0$ holds, then, for $w \in H \cap \text{cone } X$, there exists $F \in \mathcal{T}(\bar{f})$ such that $w \in F$ and $F \subseteq H$ by Lemma 2.4. Therefore, by Lemma 3.1 (2), we have $\bar{f}(w) = \overline{f^{X \cap F}}(w) = \overline{f^{X \cap H}}(w)$ and $c_H(\bar{f}) = \tilde{c}_H(f)$. \square

The theorem yields an algorithm for the split decomposition of $f : X \rightarrow \mathbf{R}$ as follows:

1. Determine \mathcal{H}_X from the points X .
2. Calculate $c_H(\bar{f})$ for $H \in \mathcal{H}_X$ by formulas (3.14) and (3.15).
3. Decompose f into the form of (3.4).

In Section 4, we derive Bandelt-Dress' split decomposition from this recipe. As will be shown in Section 4, Theorem 4.14, X -admissibility of hyperplanes corresponds to weak compatibility of splits. The next proposition corresponds to [1, Corollary 10]. This also gives a criterion for X -admissibility of hyperplanes.

Proposition 3.5. *For $\mathcal{H} \subseteq \mathcal{H}_X$ and $\alpha \in \mathbf{R}_{++}^{\mathcal{H}}$, let $f = \sum_{H \in \mathcal{H}} \alpha_H l_H^X$. Then the following conditions (a), (b) and (c) are equivalent;*

- (a) $\bar{f} = \sum_{H \in \mathcal{H}} \alpha_H l_H + \delta_{\text{cone } X}$.
- (b) $\mathcal{H} = \mathcal{H}(\bar{f})$ and $\alpha_H = c_H(\bar{f})$ for $H \in \mathcal{H}$.
- (c) \mathcal{H} is X -admissible.

Proof. (a) \Leftrightarrow (b) follows from Theorems 2.2 and 3.2. (b) \Rightarrow (c) follows from Lemma 3.3.

(c) \Rightarrow (a). Let $\mathcal{H} = \{H_{a_i, 0} \mid i \in I\}$ with $\|a_i\| = 1$ ($i \in I$). For $x \in \text{cone } X$, there exists $F \in \mathcal{A}(\mathcal{H})$ such that $x \in F$. By the condition (c), x can be represented as a conical combination

$$x = \sum_{y \in F \cap X} \lambda_y y. \quad (3.17)$$

Then the coefficient λ is feasible to the linear program corresponding (3.1). On the other hand, p defined as

$$p = \sum_{i \in I, F \subseteq H_i^+} (\alpha_{H_i}/2) a_i - \sum_{i \in I, F \subseteq H_i^-} (\alpha_{H_i}/2) a_i \quad (3.18)$$

is also feasible to the dual linear program: $\max. \langle p, x \rangle$ s.t. $p \in Q(f)$. It is easy to check that λ and p satisfy the complementary slackness condition. This implies (a). \square

In particular, from the equivalence between (b) and (c) of this proposition, we see that the decomposition into a sum of X -admissible split functions is uniquely determined. The number of X -admissible hyperplanes is bounded by $|X| - n$. This fact corresponds to [1, Corollary 4].

Proposition 3.6. *Let $\mathcal{H} \subseteq \mathcal{H}_X$ be an X -admissible set of hyperplanes. Then the set of vectors $\{l_H^X \mid H \in \mathcal{H}\} \cup \{e_i^X \mid 1 \leq i \leq n\}$ is linearly independent in \mathbf{R}^X , where $e_i : \mathbf{R}^n \rightarrow \mathbf{R}$ is the i th coordinate function defined by $e_i(x) = x_i$ for $x \in \mathbf{R}^n$. Therefore $|\mathcal{H}| \leq |X| - n$.*

Proof. Suppose that it is linearly dependent. Then there exists nonzero $(\alpha, p) \in \mathbf{R}^{\mathcal{H}} \times \mathbf{R}^n$ such that

$$\sum_{H \in \mathcal{H}, \alpha_H > 0} \alpha_H l_H^X + \langle p, \cdot \rangle = \sum_{H \in \mathcal{H}, \alpha_H < 0} -\alpha_H l_H^X. \quad (3.19)$$

By Proposition 3.5 and Lemma 3.1 (1), the convex extensions of both sides of (3.19) lead to a contradiction to the uniqueness of the polyhedral split decomposition (Theorem 2.2). \square

A function $f \in \mathbf{R}^X$ is said to be *split-decomposable* if $f - \sum_{H \in \mathcal{H}(\bar{f})} c_H(\bar{f}) l_H^X$ is (the restriction of) a linear function. By Proposition 3.6, the set of all split-decomposable functions on X can be naturally regarded as a simplicial fan of \mathbf{R}^X , which is isomorphic to the set of X -admissible sets of hyperplanes (as abstract simplicial complex). If an X -admissible \mathcal{H} has the maximal cardinality $|X| - n$, the cone

$$\left\{ \sum_{H \in \mathcal{H}} \alpha_H l_H^X + g^X \mid \alpha \in \mathbf{R}_+^{\mathcal{H}}, g : \text{linear function on } \mathbf{R}^n \right\}, \quad (3.20)$$

as a subset of \mathbf{R}^X , has interior points. Therefore, for sufficiently generic $f : X \rightarrow \mathbf{R}$ from the cone (3.20), $\mathcal{T}(f)$ forms a simplicial fan (see [14, Chapter 7]).

Corollary 3.7. *Let $\mathcal{H} \subseteq \mathcal{H}_X$ be an X -admissible set of hyperplanes with maximal cardinality $|X| - n$ and let $f = \sum_{H \in \mathcal{H}} l_H^X$. Then $\mathcal{T}^X(f)$ is a triangulation. Furthermore, the set of vertices of $\mathcal{T}^X(f)$ coincides with X .*

Proof. The latter part is immediate from the fact that if $y \in X$ is not a vertex of $\mathcal{T}^X(\bar{f})$, then \mathcal{H} is also $(X \setminus \{y\})$ -admissible. \square

In Figure 5, (b) attains the maximum number of X -admissible hyperplanes and therefore is a triangulation having vertices X .

Remark 3.8. We discuss the significance of the arguments in this section from the viewpoint of the theory of regular subdivisions and secondary fans; see [3],[14, Chapter 7] for details. Consider the equivalence relation on \mathbf{R}^X so that two functions $f, g \in \mathbf{R}^X$ are defined to be equivalent if two subdivisions of X induced by f and g coincide. Then \mathbf{R}^X is partitioned into the equivalence classes and each equivalence class is a relatively open polyhedral cone. The closure of this equivalence class is called a *secondary cone*. Let $\Sigma(X)$ be the set of secondary cones. Then $\Sigma(X)$ is a polyhedral fan and is called the *secondary fan* of X . Let $\Pi(X)$ be the fan of split-decomposable functions on X . By Propositions 3.5 and 3.6, each cone of $\Pi(X)$ is a simplicial subset of some secondary cone. Therefore, $\Pi(X)$ is a simplicial subfan of $\Sigma(X)$. See also [26, p.15] for related discussion.

Remark 3.9. We discuss the connection between the arguments in this section and Erdahl and Ryshkov's theory of *lattice dicing* [10]. For a relatively prime integer vector $z \in \mathbf{Z}^n$, we define the (infinite) set of parallel hyperplanes D_z by $\{H_{z,b} \mid b \in \mathbf{Z}\}$, where \mathbf{Z} denotes the set of integer vectors. A finite set of parallel hyperplanes (*D-family*) $\{D_{z_1}, D_{z_2}, \dots, D_{z_r}\}$ is called a *lattice dicing* with respect to \mathbf{Z}^n if the set of vertices of the subdivision $\mathcal{A}(\cup_{i=1}^r D_{z_i})$ coincides with \mathbf{Z}^n , i.e., $\cup_{i=1}^r D_{z_i}$ is \mathbf{Z}^n -admissible in our sense. A positive definite quadratic form $x^\top A x$ yields a (infinite) polyhedral subdivision to \mathbf{R}^n by projecting the lower faces of the convex hull of $\{(x, x^\top A x) \mid x \in \mathbf{Z}^n\} \subseteq \mathbf{R}^n \times \mathbf{R}$, which is called an *L-partition*. Analogously to the case of secondary fans (Remark 3.8), we define the equivalence relation on the cone of positive definite quadratic forms so

that two positive definite forms are defined to be equivalent if the corresponding L -partitions coincide. Then the cone of positive definite quadratic forms is partitioned into the equivalence classes called L -type domains, which are known to be relatively open polyhedral cones. The domain of lattice dicing $\{D_{z_1}, D_{z_2}, \dots, D_{z_r}\}$ is a set of positively definite forms represented by $\sum_{i=1}^r \lambda_i (\langle d_i, x \rangle)^2$ for positive weights $\{\lambda_i\}_{i=1}^r$. Erdahl and Ryshkov [10] show that L -partition induced by $\sum_{i=1}^r \lambda_i (\langle d_i, x \rangle)^2$ coincides with $\mathcal{A}(\cup_{i=1}^r D_{z_i})$, and each dicing domain is a simplicial L -type domain. So the fan of (the closure of) dicing domains can be understood as a correspondence of the fan of split-decomposable functions. Indeed, quadratic form $(\langle d_i, x \rangle)^2$ on $\{x \in \mathbf{Z}^n \mid -m \leq \langle d_i, x \rangle \leq m\}$ can be represented as (the restriction of) a weighted sum of split functions $\{l_{H_{d_i, k}}\}_{k=-m+1}^{m-1}$ and an affine function.

4 Metrics as discrete concave functions

By regarding metrics as discrete concave functions and applying the results of Sections 2 and 3, we derive in this section Bandelt-Dress' split decomposition of metrics and some other important concepts of T-theory.

First we briefly review T-theory and the split decomposition of metrics. Let V be a finite set. A function $d : V \times V \rightarrow \mathbf{R}$ is said to be a *metric* if it satisfies $d(i, i) = 0$, $d(i, j) = d(j, i) \geq 0$, and $d(i, j) \leq d(i, k) + d(j, k)$ for $i, j, k \in V$. A polyhedron $P(d) \subseteq \mathbf{R}^V$ associated with metric d is defined as

$$P(d) = \{p \in \mathbf{R}^V \mid p(i) + p(j) \geq d(i, j)\}. \quad (4.1)$$

The tight span of metric d is a subset of $P(d)$ defined as

$$T(d) = \{p \in \mathbf{R}^V \mid \forall i \in V, p(i) = \sup_{j \in V} (d(i, j) - p(j))\}. \quad (4.2)$$

It is known that $T(d)$ coincides with the union of all bounded faces of $P(d)$ [8, Lemma 1].

A *split* $\{A, B\}$ is a bipartition of V , which means that $A \cap B = \emptyset$, $A \cup B = V$, $A, B \neq \emptyset$. A *split metric* $\delta_{\{A, B\}}$ associated with a split $\{A, B\}$ is defined as

$$\delta_{\{A, B\}}(i, j) = \begin{cases} 0 & \text{if } i, j \in A \text{ or } i, j \in B, \\ 1 & \text{otherwise} \end{cases} \quad (i, j \in V). \quad (4.3)$$

For a metric d and a split $\{A, B\}$, the *isolation index* $\alpha_{\{A, B\}}^d$ is the nonnegative number defined as

$$\alpha_{\{A, B\}}^d = \frac{1}{2} \min_{i, j \in A, k, l \in B} \left\{ \max \begin{cases} d(i, k) + d(j, l) \\ d(i, l) + d(j, k) \\ d(i, j) + d(k, l) \end{cases} - d(i, j) - d(k, l) \right\}. \quad (4.4)$$

Let $\mathcal{S}(d)$ be the collection of splits defined as

$$\mathcal{S}(d) = \{S : \text{split on } V : \alpha_S^d > 0\}. \quad (4.5)$$

An additive decomposition of a metric $d = d_1 + d_2$ is said to be *coherent* if it satisfies $P(d) = P(d_1) + P(d_2)$. A metric d is said to be *split-prime* if it satisfies $\alpha_S^d = 0$ for each split S . The split decomposition theorem is as follows:

Theorem 4.1 (Bandelt-Dress [1]). *A metric d can be coherently decomposed as*

$$d = \sum_{S \in \mathcal{S}(d)} \alpha_S^d \delta_S + d', \quad (4.6)$$

where d' is a split-prime metric. Moreover, if d is coherently decomposed as

$$d = \sum_{S \in \mathcal{S}} \lambda_S \delta_S + d'' \quad (4.7)$$

for a collection of splits \mathcal{S} , a positive weight $\lambda \in \mathbf{R}_{++}^S$, and a split-prime metric d'' . then we have $\mathcal{S} = \mathcal{S}(d)$, $\lambda_S = \alpha_S^d$ for $S \in \mathcal{S}$, and $d'' = d'$.

The collection of splits $\mathcal{S}(d)$ is necessarily *weakly compatible*, that is, for any three splits $S_1 = \{A_1, B_1\}$, $S_2 = \{A_2, B_2\}$, and $S_3 = \{A_3, B_3\}$ in $\mathcal{S}(d)$, there exist no four points $a, a_1, a_2, a_3 \in V$ with $\{a, a_1, a_2, a_3\} \cap A_i = \{a, a_i\}$ for $i = 1, 2, 3$.

As already mentioned in the introduction, we regard a metric as a function defined on a finite set of points. In the following we deal with a more general object, *distance function*, rather than a metric. A distance function on V is a function $d : V \times V \rightarrow \mathbf{R}$ such that $d(i, i) = 0$ and $d(i, j) = d(j, i) \geq 0$ for $i, j \in V$, where the triangle inequality is not imposed. For each $A \subseteq V$, we denote by χ_A the characteristic vector of A defined as: $\chi_A(i) = 1$ if $i \in A$ and 0 otherwise. In particular, we write simply χ_i instead of $\chi_{\{i\}}$ for each $i \in V$. Consider the point set

$$\Lambda = \{\chi_i + \chi_j \mid i, j \in V\}. \quad (4.8)$$

Then a distance function d is naturally regarded as a function $d : \Lambda \rightarrow \mathbf{R}$ defined by

$$d(\chi_i + \chi_j) = d(i, j) \quad (i, j \in V). \quad (4.9)$$

The following is easily observed.

Lemma 4.2. *A function $f : \Lambda \rightarrow \mathbf{R}$ with $f(2\chi_i) = 0$ for $i \in V$ is convex-extensible if and only if it satisfies $f(\chi_i + \chi_j) \leq 0$ for $i, j \in V$.*

Hence it is natural to regard any distance function $d : V \times V \rightarrow \mathbf{R}$ as a discrete concave function on Λ . Since $\text{aff } \Lambda$ coincides with $K = \{x \in \mathbf{R}^V \mid \sum_{i \in V} x(i) = 2\}$, we can apply the results of the previous section in a straightforward manner.

The homogeneous convex extension of $-d$ is given by

$$\begin{aligned} & \overline{(-d)}(x) \\ &= \inf \left\{ \sum_{i, j \in V} \lambda_{ij} (-d(i, j)) \mid \sum_{i, j \in V} \lambda_{ij} (\chi_i + \chi_j) = x, \lambda_{ij} \geq 0 \quad (i, j \in V) \right\} \\ &= \sup \{ \langle p, x \rangle \mid p(i) + p(j) \leq -d(i, j) \} \\ &= \sup \{ \langle p, x \rangle \mid -p \in P(d) \} \quad (x \in \mathbf{R}^V), \end{aligned} \quad (4.10)$$

where for $i, j \in V$, ij denotes an unordered pair, which means that ij and ji are not distinguished from each other. Hence $\overline{-d}$ is the support function of $-P(d)$. Using the notation (3.2), we have $Q(-d) = -P(d)$. This implies that the tight span $T(d)$ has a dual structure of $\mathcal{T}^\Lambda(\overline{-d})$ (recall (N1), (N2), and (N3) in Subsection 2.3 and see Figure 6). This duality relation is also suggested by Sturmfels and Yu [26] in the connection of the triangulation of the *second hypersimplex* $\text{conv}\{\chi_i + \chi_j \mid i, j \in V, i \neq j\}$.

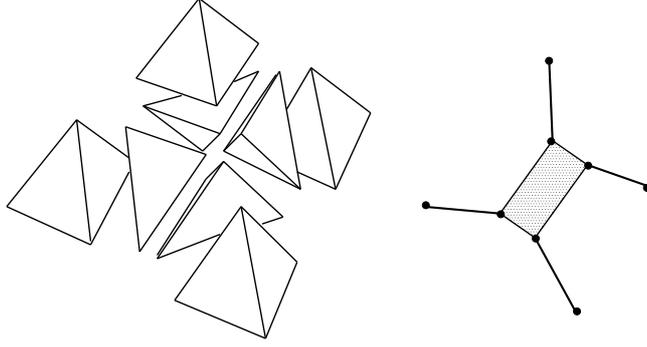


Figure 6: $T^\Lambda(\overline{-d})$ (left) and $T(d)$ (right) of a generic 4-point metric d

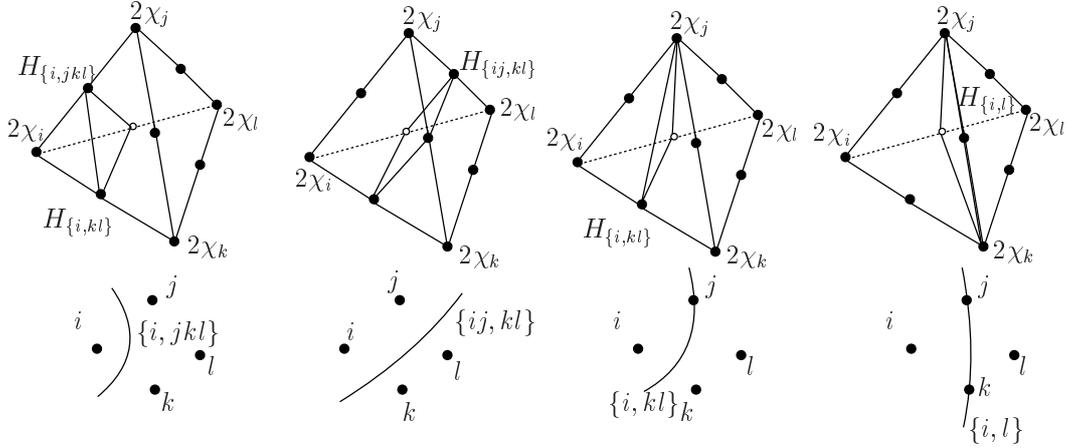


Figure 7: Examples of the split-hyperplane correspondence

Next we derive the split decomposition of metrics (Theorem 4.1) in a more generalized form using the result of the present paper. For this, we begin by establishing a relationship between splits and hyperplanes. For $A, B \subseteq V$ with $A, B \neq \emptyset$ and $A \cap B = \emptyset$, we call unordered pair $\{A, B\}$ a *partial split*. A hyperplane $H_{\{A, B\}}$ associated with partial split $\{A, B\}$ is defined as

$$H_{\{A, B\}} = \{x \in \mathbf{R}^V \mid x(A) = x(B)\}, \quad (4.11)$$

where we denote $\sum_{i \in A} x(i)$ by $x(A)$ for $x \in \mathbf{R}^V$ and $A \subseteq V$.

Lemma 4.3. *Let $\{A, B\}$ be a partial split on V and $C = V \setminus (A \cup B)$. Then we have*

$$\begin{cases} \Lambda \cap H_{\{A, B\}} &= \{\chi_i + \chi_j \mid i \in A, j \in B \text{ or } i, j \in C\}, \\ \Lambda \cap H_{\{A, B\}}^+ &= \{\chi_i + \chi_j \mid i, j \in A \cup C\}, \\ \Lambda \cap H_{\{A, B\}}^- &= \{\chi_i + \chi_j \mid i, j \in B \cup C\}, \\ \Lambda \cap H_{\{A, B\}}^{++} &= \{\chi_i + \chi_j \mid i \in A, j \in A \cup C\}, \\ \Lambda \cap H_{\{A, B\}}^{--} &= \{\chi_i + \chi_j \mid i \in B, j \in A \cup C\}, \end{cases} \quad (4.12)$$

where we regard $H_{\{A, B\}}$ as $H_{\chi_A - \chi_B, 0}$.

Figure 7 illustrates small examples of the correspondence between partial splits and hyperplanes. In this figure, a partial split and an element of V are represented by a line and a point, respectively. The line corresponding to a partial split $\{A, B\}$ separates

points of A and B and contains points of $V \setminus A \cup B$. The family \mathcal{H}_Λ of linear hyperplanes, as defined in (3.13) together with (a0) and (a1), can be identified in terms of partial splits as follows.

Proposition 4.4. $\mathcal{H}_\Lambda = \{H_{\{A,B\}} \mid \{A,B\} \text{ is a partial split on } V\}$.

Proof. We show that $H_{A,B}$ satisfies the conditions (a0) and (a1). (a0) is clearly satisfied. $\text{conv } \Lambda$ is the simplex having edges $\{[2\chi_i, 2\chi_j] \mid i, j \in V\}$. By Lemma 4.3, $2\chi_i \in H_{A,B}^{++}$ and $2\chi_j \in H_{A,B}^{--}$ if and only if $i \in A$ and $j \in B$. Therefore, we have $[2\chi_i, 2\chi_j] \cap H_{A,B} = \{\chi_i + \chi_j\} \in X$ for $i \in A, j \in B$. Conversely, let $H_{a,0} \in \mathcal{H}_\Lambda$. Define $A = \{i \in V \mid 2\chi_i \in H_{a,0}^{++}\}$ and $B = \{i \in V \mid 2\chi_i \in H_{a,0}^{--}\}$. Then $\{A, B\}$ is a partial split on V . We show $H_{a,0} = H_{\{A,B\}}$. By the definition of $\{A, B\}$, we have $a(i) = 0$ for $i \in V \setminus (A \cup B)$. For any $2\chi_i \in H_{a,0}^{++}$ and $2\chi_j \in H_{a,0}^{--}$, by Λ -admissibility of $\{H_{a,0}\}$, it is necessary that $\chi_i + \chi_j \in H_{a,0}$ by (a1). Hence we have $a(i) = -a(j)$ for $i \in A, j \in B$. This implies $H_{a,0} = H_{\{A,B\}}$. \square

For a partial split $\{A, B\}$, we define a *partial split distance* $\zeta_{\{A,B\}}$ associated with $\{A, B\}$ as

$$\zeta_{\{A,B\}}(i, j) = \begin{cases} 1 & \text{if } i \in A, j \in B \text{ or } i \in B, j \in A \\ 0 & \text{otherwise} \end{cases} \quad (i, j \in V). \quad (4.13)$$

If $A \cup B \neq V$, then $\zeta_{\{A,B\}}$ is not a metric. If $A \cup B = V$, i.e., $\{A, B\}$ is a split, then $\zeta_{\{A,B\}}$ coincides with the split metric $\delta_{\{A,B\}}$.

The following shows that a partial split distance $\zeta_{\{A,B\}}$ associated with partial split $\{A, B\}$ is represented as split function $l_{H_{\{A,B\}}}$ associated with hyperplane $H_{\{A,B\}}$.

Proposition 4.5. *For a partial split $\{A, B\}$, the partial split distance $\zeta_{\{A,B\}} : \Lambda \rightarrow \mathbf{R}$ is represented as*

$$\zeta_{\{A,B\}}(x) = -|x(A) - x(B)|/2 + x(A \cup B)/2 \quad (x \in \Lambda). \quad (4.14)$$

Moreover, we have

$$\overline{(-\zeta_{\{A,B\}})}(x) = |x(A) - x(B)|/2 - x(A \cup B)/2 + \delta_{\mathbf{R}_+^V} \quad (x \in \mathbf{R}^V), \quad (4.15)$$

where $\delta_{\mathbf{R}_+^V}$ is the indicator function for \mathbf{R}_+^V ; see (2.5).

Proof. The equation (4.14) is obtained by direct calculations. The equation (4.15) follows from Λ -admissibility of $\{H_{\{A,B\}}\}$ (Proposition 4.4), Proposition 3.5, and Lemma 3.1 (1). \square

$c_{H_{\{A,B\}}}(-d)$ can be explicitly calculated as follows.

Theorem 4.6. *For a distance function d and a partial split $\{A, B\}$, we have*

$$c_{H_{\{A,B\}}}(\overline{(-d)}) = \max\{0, \min\{\rho_{A,B}^d, \sigma_{A,B}^d, \sigma_{B,A}^d, \tau_{A,B}^d\}\} \sqrt{|A \cup B|}, \quad (4.16)$$

where

$$\rho_{A,B}^d = \frac{1}{2} \min \left\{ \max \left\{ \begin{array}{l} d(i, k) + d(j, l) \\ d(i, l) + d(j, k) \end{array} \right\} - d(i, j) - d(k, l) \mid \begin{array}{l} i, j \in A \\ k, l \in B \end{array} \right\}, \quad (4.17)$$

$$\sigma_{A,B}^d = \frac{1}{2} \min \left\{ d(i, k) + d(i, l) - d(k, l) - 2d(i, j) \mid \begin{array}{l} i \in A, k, l \in B \\ j \in V \setminus (A \cup B) \end{array} \right\}, \quad (4.18)$$

$$\tau_{A,B}^d = \min \left\{ d(i, k) + d(j, l) - d(i, j) - d(k, l) \mid \begin{array}{l} i \in A, k \in B \\ j, l \in V \setminus (A \cup B) \end{array} \right\} \quad (4.19)$$

and we define $\sigma_{A,B}^d = +\infty$ and $\tau_{A,B}^d = +\infty$ if $A \cup B = V$.

Proof. We apply the formulas (3.14) and (3.15) in Theorem 3.4. Let $C = V \setminus (A \cup B)$. Then $\tilde{c}_{H_S}(-d)$ is the minimum of

$$\frac{-d(i, j) - \overline{(-d)^{\Lambda \cap H_{\{A, B\}}}}(w)}{2l_{H_{\{A, B\}}}(\chi_i + \chi_j)} + \frac{-d(k, l) - \overline{(-d)^{\Lambda \cap H_{\{A, B\}}}}(w)}{2l_{H_{\{A, B\}}}(\chi_k + \chi_l)} \quad (4.20)$$

over $i \in A, j \in A \cup C, k \in B, \text{ and } l \in B \cup C$, where $\{w\} = H_{\{A, B\}} \cap [\chi_i + \chi_j, \chi_k + \chi_l]$.
(Case 1: $i, j \in A, k, l \in B$). (4.20) is given by

$$\frac{\sqrt{|A \cup B|}}{2} \left\{ -d(i, j) - d(k, l) - 2\overline{(-d)^{\Lambda \cap H_S}}((\chi_i + \chi_j + \chi_k + \chi_l)/2) \right\}.$$

Furthermore $\overline{(-d)^{\Lambda \cap H_{\{A, B\}}}}((\chi_i + \chi_j + \chi_k + \chi_l)/2)$ is given by the optimal value of linear program

$$\begin{aligned} \min. \quad & \sum_{uv \in \{ik, il, jk, jl\}} \lambda_{uv}(-d(u, v)) \\ \text{s.t.} \quad & \sum_{uv \in \{ik, il, jk, jl\}} \lambda_{uv}(\chi_u + \chi_v) = (\chi_i + \chi_j + \chi_k + \chi_l)/2, \\ & \lambda_{uv} \geq 0 \quad (uv \in \{ik, il, jk, jl\}). \end{aligned}$$

By direct calculations, the optimal value of the problem is

$$\min\{-d(i, k) - d(j, l), -d(i, l) - d(j, k)\}/2.$$

Hence we have

$$(4.20) = \frac{\sqrt{|A \cup B|}}{2} \left\{ -d(i, j) - d(k, l) + \max \left\{ \begin{array}{l} d(i, k) + d(j, l) \\ d(i, l) + d(j, k) \end{array} \right\} \right\}. \quad (4.21)$$

(Case 2: $i \in A, j \in C, k, l \in B$). (4.20) is given by

$$\frac{\sqrt{|A \cup B|}}{2} \left\{ -2d(i, j) - d(k, l) - 3\overline{(-d)^{\Lambda \cap H_{\{A, B\}}}}((2\chi_i + 2\chi_j + \chi_k + \chi_l)/3) \right\}.$$

In a similar way of (Case 1), we have

$$\overline{(-d)^{\Lambda \cap H_{\{A, B\}}}}((2\chi_i + 2\chi_j + \chi_k + \chi_l)/3) = -(d(i, k) + d(i, l))/3.$$

Hence we obtain

$$(4.20) = \frac{\sqrt{|A \cup B|}}{2} \{-2d(i, j) - d(k, l) + d(i, k) + d(i, l)\}. \quad (4.22)$$

(Case 3: $i, j \in A, k \in B, l \in C$). It follows from (Case 2) by interchanging A and B .
(Case 4: $i \in A, k \in B, j, l \in C$). (4.20) is given by

$$\sqrt{|A \cup B|} \left\{ -d(i, j) - d(k, l) - 2\overline{(-d)^{\Lambda \cap H_{\{A, B\}}}}((\chi_i + \chi_j + \chi_k + \chi_l)/2) \right\}.$$

In a similar way of (Case 1), we have

$$\overline{(-d)^{\Lambda \cap H_{\{A, B\}}}}((\chi_i + \chi_j + \chi_k + \chi_l)/2) = -(d(i, k) + d(j, l))/2.$$

Hence, we obtain

$$(4.20) = \sqrt{|A \cup B|} \{-d(i, j) - d(k, l) + d(i, k) + d(j, l)\}. \quad (4.23)$$

Combining (4.21), (4.22), and (4.23), we obtain the desired formula (4.16). \square

In particular, if a partial split $\{A, B\}$ forms a split, then we have

$$c_{H_{\{A,B\}}}(\overline{-d}) = \alpha_{\{A,B\}}^d \sqrt{|V|} \quad (4.24)$$

(see the definitions of $\alpha_{\{A,B\}}^d$ in (4.4) and $\rho_{A,B}^d$ in (4.17)). Accordingly, the isolation index can be extended for a partial split $\{A, B\}$ by defining

$$\alpha_{\{A,B\}}^d = \max\{0, \min\{\rho_{A,B}^d, \sigma_{A,B}^d, \sigma_{B,A}^d, \tau_{A,B}^d\}\}. \quad (4.25)$$

Similarly, $\mathcal{S}(d)$ can be extended for partial splits by defining

$$\mathcal{S}(d) = \{S : \text{partial split on } V : \alpha_S^d > 0\}. \quad (4.26)$$

A distance d is said to be *split-prime* if $\alpha_S^d = 0$ for each partial split S . An additive decomposition of a distance $d = d_1 + d_2$ is said to be *coherent* if $P(d) = P(d_1) + P(d_2)$. Note that this condition is equivalent to $\overline{(-d)} = \overline{(-d_1)} + \overline{(-d_2)}$.

As a consequence of above arguments, we obtain an extension of Bandelt-Dress' split decomposition (Theorem 4.1).

Theorem 4.7. *A distance function $d : V \times V \rightarrow \mathbf{R}$ can be coherently decomposed as*

$$d = \sum_{S \in \mathcal{S}(d)} \alpha_S^d \zeta_S + d', \quad (4.27)$$

where d' is a split-prime distance. Moreover, if d is coherently decomposed as

$$d = \sum_{S \in \mathcal{S}} \lambda_S \zeta_S + d'' \quad (4.28)$$

for a collection of partial splits \mathcal{S} , a positive weight $\lambda \in \mathbf{R}_{++}^{\mathcal{S}}$ and a split-prime distance d'' , then we have $\mathcal{S} = \mathcal{S}(d)$, $\lambda_S = \alpha_S^d$ for $S \in \mathcal{S}(d)$, and $d'' = d'$.

Proof. Decompose $-d$ into the form (3.4) in Theorem 3.2 and apply Propositions 4.4 and 4.5, Theorem 4.6 and Lemma 4.2. Then we obtain (4.27). The coherency of the decomposition (4.27) follows from $Q(-d) = -P(d)$. We show the latter part. By $\sum_{S \in \mathcal{S}(d)} \alpha_S^d \overline{(-\zeta_S)} + \overline{(-d')} = \sum_{S \in \mathcal{S}} \lambda_S \overline{(-\zeta_S)} + \overline{(-d'')}$, Proposition 4.5, and Theorem 2.2. We have $\mathcal{S}(d) = \mathcal{S}$, $\alpha_S^d = \lambda_S$ and $\overline{(-d')} = \overline{(-d'')}$. By convex-extensibility of distances (Lemma 4.2), we have $-d' = \overline{(-d')}^\Lambda = \overline{(-d'')}^\Lambda = -d''$. \square

We call (4.27) the *extended split decomposition* of d . The cardinality of $\mathcal{S}(d)$ is bounded by $|V|(|V|-1)/2$ by Proposition 3.6. Therefore the decomposition (4.27) is also obtained in polynomial time by an algorithm similar to that for the split decomposition of a metric [1]. It is easily observed that if d satisfies the triangle inequality, $\tau_{\{A,B\}}^d \leq 0$ holds for any proper partial split $\{A, B\}$ with $A \cup B \neq V$. This implies that $\mathcal{S}(d)$ consists of splits. Hence if d is a metric, the decomposition (4.27) coincides with the split decomposition of metrics (4.6).

Remark 4.8. Any symmetric function $f : V \times V \rightarrow \mathbf{R}$ with $f(i, j) = f(j, i)$ for $i, j \in V$ can be regarded as $f : \Lambda \rightarrow \mathbf{R}$ by the correspondence (4.9). Similarly, $P(f)$ and $T(f)$ are also definable. The arguments in this section can be adapted to any symmetric function on V . In fact, [1] discusses the split decomposition (4.6) for symmetric functions. However, [1] uses only split metrics in the decomposition (4.6). In the case that d violates triangle inequality, it is possible to achieve a more precise decomposition of d by (4.27)

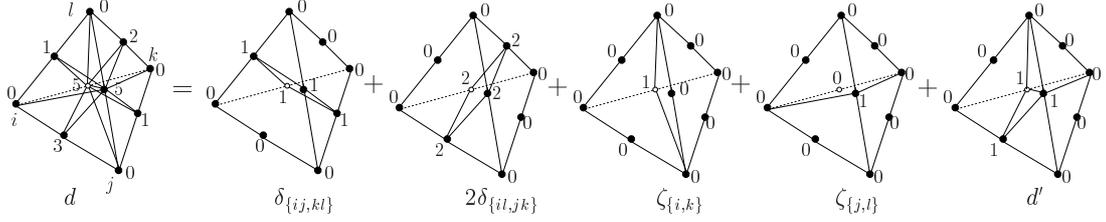


Figure 8: An example of the extended split decomposition

since (4.27) uses also partial split distances. For example, consider the following 4-point distance function

$$d = \begin{array}{c|cccc} & i & j & k & l \\ \hline i & 0 & 3 & 5 & 1 \\ j & 3 & 0 & 1 & 5 \\ k & 5 & 1 & 0 & 2 \\ l & 1 & 5 & 2 & 0 \end{array} . \quad (4.29)$$

Note that d violates the triangle inequality. By the extended split decomposition (4.27), d is decomposed as

$$d = \delta_{\{ij,kl\}} + 2\delta_{\{ik,jl\}} + \zeta_{\{i,k\}} + \zeta_{\{j,l\}} + \begin{array}{c|cccc} & i & j & k & l \\ \hline i & 0 & 1 & 1 & 0 \\ j & 1 & 0 & 0 & 1 \\ k & 1 & 0 & 0 & 0 \\ l & 0 & 1 & 0 & 0 \end{array} \quad (4.30)$$

In the ordinary Bandelt-Dress' split decomposition (4.6), only the first two term of RHS of (4.30) are exploited and the sum of the last three terms is treated as the split-prime residue. Figure 8 illustrates this decomposition in term of the subdivision $\mathcal{T}^\Lambda(\overline{-d})$. This indicates that the extended split decomposition exploits more detailed combinatorial structure of distances than Bandelt-Dress' split decomposition.

Remark 4.9. It is well known that every 4-point metric is totally split-decomposable, i.e., a split-prime residue d' of (4.6) vanishes. It is easily seen that every 3-point distance function is also totally split-decomposable in our sense, i.e., $d' = 0$ in the decomposition (4.27). However, as is seen in Remark 4.8, not every 4-point distance function is totally split-decomposable in our sence.

Remark 4.10. By the decomposition (4.27), $P(d)$ is decomposed as

$$P(d) = Z(d) + P(d'), \quad (4.31)$$

where $Z(d)$ is (a translation of) the maximum zonotopic summand given as

$$Z(d) = \sum_{\{A,B\} \in \mathcal{S}(d)} \alpha_{\{A,B\}}^d ([\chi_A - \chi_B, \chi_B - \chi_A]/2 + \chi_{A \cup B}/2) \quad (4.32)$$

If $d' = 0$. we have $P(d) = Z(d) + \mathbf{R}_+^V$. In this case, tight span $T(d)$ is the union of the faces of $Z(d)$ whose normal cone contains negative vectors.

Remark 4.11. The coherent decomposition decomposes a distance d into a sum of distances which induce a coarser subdivision than $\mathcal{T}^\Lambda(\overline{-d})$. In particular, it is observed

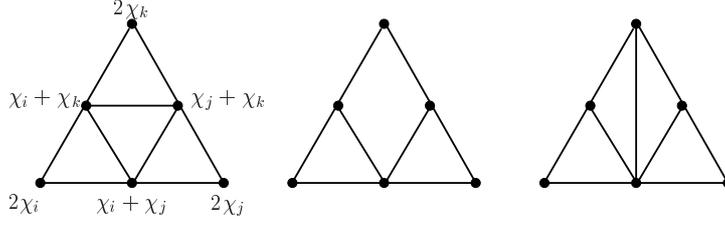


Figure 9: The role of triangle inequality: $d(i, j) < d(i, k) + d(j, k)$ (left), $d(i, j) = d(i, k) + d(j, k)$ (center) and $d(i, j) > d(i, k) + d(j, k)$ (right)

from the subdivision of $\text{conv}\{2\chi_i, 2\chi_j, 2\chi_k\}$ induced by $\mathcal{T}^\Lambda(\overline{(-d)})$ that d satisfies $d(i, j) \leq d(i, k) + d(j, k)$ if and only if $\text{conv}\{2\chi_k, \chi_i + \chi_j\} \notin \mathcal{T}^\Lambda(\overline{(-d)})$ (see Figure 9). This implies that in any coherent decomposition $d = d_1 + d_2$, if d is a metric, then both d_1 and d_2 are necessarily metrics.

Remark 4.12. Koolen, Moulton and Tönges [18] introduced the *coherency index* as a direct generalization of isolation index. For two metrics d' and d , the coherency index is the nonnegative value $\alpha_{d'}^d$ which has the property that $d = \alpha d' + (d - \alpha d')$ is a coherent decomposition if and only if $0 \leq \alpha \leq \alpha_{d'}^d$ [18, Theorem 4.1]. By definition, $\alpha_{\zeta_S}^d = \alpha_S^d$ must hold for a partial split S . They give a formula of $\alpha_{d'}^d$ as

$$\alpha_{d'}^d = \min_{f: \text{vertex of } P(d)} \max_{f': \text{vertex of } P(d')} m(f, f'), \quad (4.33)$$

where $m(f, f')$ is defined as

$$m(f, f') = \min_{i, j \in V: f'(i) + f'(j) \neq d'(i, j)} \frac{f(i) + f(j) - d(i, j)}{f'(i) + f'(j) - d'(i, j)}. \quad (4.34)$$

This formula come from the fact that $P(d) = \alpha P(d') + P(d - \alpha d')$ if and only if each $f \in P(d)$ and $f' \in P(d')$ satisfies $f - \alpha f' \in P(d - \alpha d')$. It seems to be difficult to derive the explicit formula (4.25) of $\alpha_{\zeta_S}^d = \alpha_S^d$ from the coherency index (4.33). It should be noted that our derivation of α_S^d come from the fact that $P(d) = \alpha P(d') + P(d - \alpha d')$ if and only if $\overline{(-d)} - \alpha \overline{(-d')}$ is convex on \mathbf{R}_+^V .

Next we characterize Λ -admissibility of hyperplanes in terms of combinatorial properties of the corresponding partial splits. First we observe that $\mathcal{S}(d)$ does not contain the following types of partial splits (see Figure 10):

- (C1) three partial splits $\{A_1, B_1\}$, $\{A_2, B_2\}$, $\{A_3, B_3\}$ and four points $a, a_1, a_2, a_3 \in V$ such that $\{a, a_1, a_2, a_3\} \cap A_i = \{a, a_i\}$ for $i = 1, 2, 3$ (the violation of weak compatibility).
- (C2) two partial splits $\{A_1, B_1\}$, $\{A_2, B_2\}$ and three points $a, b, c \in V$ such that $a \in A_1$, $b, c \in B_1$, $b \in A_2$, $c \in B_2$, and $a \in V \setminus (A_2 \cup B_2)$.
- (C3) two partial splits $\{A_1, B_1\}$, $\{A_2, B_2\}$ and three points $a, b, c \in V$ such that $b \in A_1$, $c \in B_1$, $a \in V \setminus (A_1 \cup B_1)$, $a \in A_2$, $b \in B_2$, and $c \in V \setminus (A_2 \cup B_2)$.

For (C1), observe that not all three ρ_{A_1, B_1}^d , ρ_{A_2, B_2}^d and ρ_{A_3, B_3}^d are positive. Similarly, for (C2), observe that both ρ_{A_1, B_1}^d and τ_{A_2, B_2}^d are not positive. To see (C3), observe that

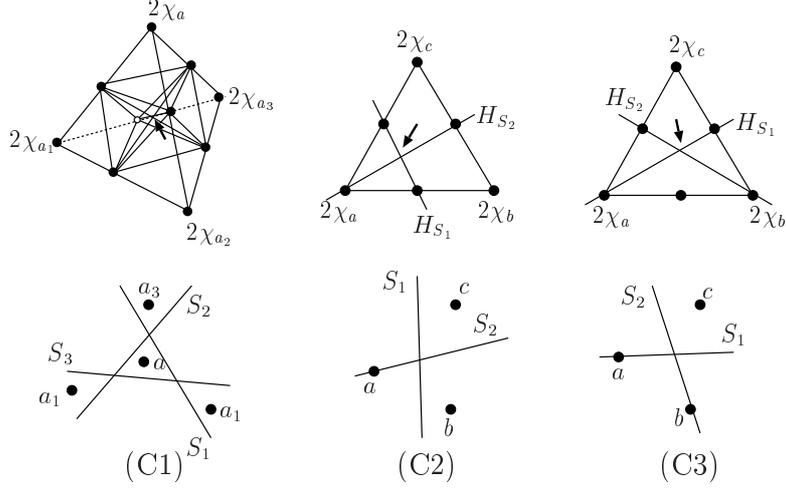


Figure 10: Forbidden partial splits and the corresponding hyperplanes

both τ_{A_1, B_1}^d and τ_{A_2, B_2}^d are not positive. Of course, these can be geometrically checked by the corresponding hyperplanes in Figure 10; in each case, there is a vertex, indicated by the arrow, not contained by Λ .

A collection of partial splits free from (C1), (C2) and (C3) is said to be *weakly compatible*, which is clearly an extension of weakly compatibility of splits. The following is an extension of [1, Theorem 3].

Theorem 4.13. *Let \mathcal{S} be a weakly compatible collection of partial splits on V . For $\lambda \in \mathbf{R}_{++}^S$, let a distance function $d : V \times V \rightarrow \mathbf{R}$ be defined as $d = \sum_{S \in \mathcal{S}} \lambda_S \zeta_S$. Then we have $\mathcal{S} = \mathcal{S}(d)$ and $\lambda_S = \alpha_S^d$ for each $S \in \mathcal{S}$.*

Proof. We adapt the proof of [1, Theorem 3]. For a partial split $\{A, B\}$, let $\tilde{\alpha}_{\{A, B\}}^d$ be defined as

$$\tilde{\alpha}_{\{A, B\}}^d = \min\{\rho_{A, B}^d, \sigma_{A, B}^d, \sigma_{B, A}^d, \tau_{A, B}^d\}.$$

It suffices to show

$$\tilde{\alpha}_S^d \geq \lambda_S \quad (S \in \mathcal{S}). \quad (4.35)$$

By the formula of α_S^d for a partial split $S = \{A, B\}$, there exists $Y \subseteq V$ with $A \cap Y \neq \emptyset$, $B \cap Y \neq \emptyset$, and $|Y| \leq 4$ such that

$$\tilde{\alpha}_{\{A, B\}}^d = \tilde{\alpha}_{\{A \cap Y, B \cap Y\}}^{d_Y},$$

where $d_Y : Y \times Y \rightarrow \mathbf{R}$ denotes the restriction of d to Y . Hence it is sufficient to show (4.35) for the case $|V| \leq 4$.

In the case of $|V| = 2$, (4.35) is obvious. In the case of $|V| = 3$, let $V = \{i, j, k\}$. For the simplicity of notation, we denote a partial split $\{\{i, j\}, \{k\}\}$ by $\{ij, k\}$. It suffices to check (4.35) for $S = \{ij, k\}$ and $S = \{i, j\}$. For $S = \{ij, k\}$, \mathcal{S} does not contain $\{i, j\}$ by (C2). Hence we have $\tilde{\alpha}_{\{ij, k\}}^d = \rho_{ij, k}^d = \min\{(d(i, k) + d(j, k) - d(i, j))/2, d(i, k), d(j, k)\} \geq \lambda_{\{ij, k\}}$. For $S = \{i, j\}$, \mathcal{S} does not contain $\{j, k\}$, $\{i, k\}$ and $\{ij, k\}$ from (C2) and (C3). We have $\tau_{i, j}^d = d(i, j) - d(i, k) - d(j, k) = \lambda_{\{i, j\}}$, $\rho_{i, j}^d = d(i, j) \geq \lambda_{\{i, j\}}$, $\sigma_{i, j}^d \geq \tau_{i, j}^d$ and $\sigma_{j, i}^d \geq \tau_{i, j}^d$. Hence we obtain $\tilde{\alpha}_{\{i, j\}}^d \geq \lambda_{\{i, j\}}$. Therefore (4.35) holds for $|V| = 3$.

In the case of $|V| = 4$, let $V = \{i, j, k, l\}$. We may assume that the minimum of $\tilde{\alpha}_S^d$ is attained by all different four points (otherwise it can be reduced to the case $|V| \leq 3$).

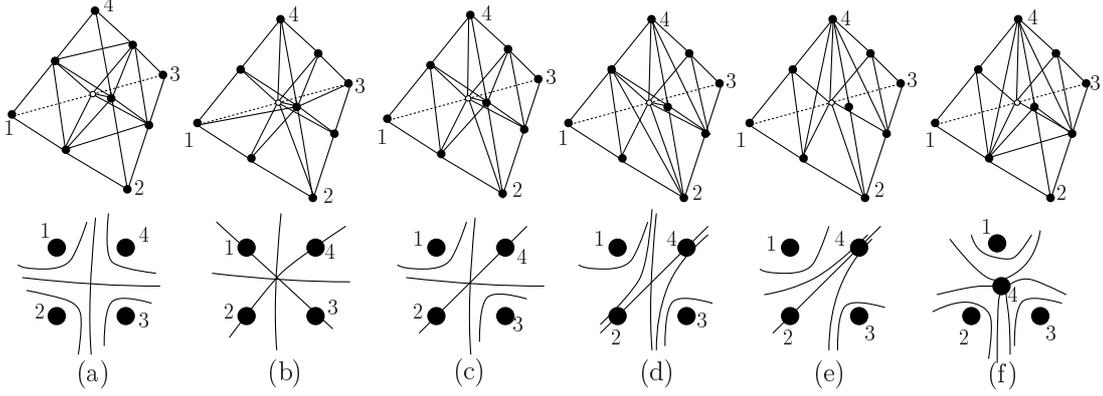


Figure 11: Maximal weakly compatible partial splits on 4-point set and the corresponding hyperplanes

It suffices to check (4.35) for $S = \{ij, kl\}$, $S = \{ij, k\}$ and $S = \{i, j\}$. For $S = \{ij, kl\}$, we have

$$\tilde{\alpha}_{\{ij,kl\}}^d = \rho_{ij,kl}^d = \{\max\{d(i, k) + d(j, l), d(i, l) + d(j, k)\} - d(i, j) - d(k, l)\}/2. \quad (4.36)$$

From the condition (C1), \mathcal{S} does not contain both $\{ik, jl\}$ and $\{il, jk\}$ simultaneously. Suppose that $\{ik, jl\} \in \mathcal{S}$ and $\{il, jk\} \notin \mathcal{S}$. By the conditions (C2) and (C3), the possible partial splits in \mathcal{S} that contribute to the term $d(i, j) + d(k, l)$ are only five partial splits $\{i, jkl\}$, $\{j, ikl\}$, $\{k, ijl\}$, $\{l, ij k\}$, and $\{ik, jl\}$. First four partial splits contribute equally to $d(i, k) + d(j, l)$, $d(i, l) + d(j, k)$ and $d(i, j) + d(k, l)$. Hence we have $\tilde{\alpha}_{\{ij,kl\}}^d \geq \max\{\lambda_{\{ij,kl\}}, \lambda_{\{ij,kl\}} + \lambda_{\{ik,jl\}}\} - \lambda_{\{ik,jl\}} = \lambda_{\{ij,kl\}}$. For $S = \{ij, k\}$, we have

$$\tilde{\alpha}_{\{ij,k\}}^d = \sigma_{k,ij}^d = \{d(i, k) + d(j, k) - d(i, j) - 2d(k, l)\}/2. \quad (4.37)$$

From the conditions (C2) and (C3), the possible partial splits in \mathcal{S} that contribute to the term $d(i, j) + 2d(k, l)$ are four partial splits $\{i, jkl\}$, $\{j, ikl\}$, $\{k, ijl\}$, and $\{ik, j\}$. However it is easily examined that these partial split distances cancel out in (4.37). Therefore we obtain $\tilde{\alpha}_{\{ij,k\}}^d \geq \lambda_{\{ij,k\}}$. For $S = \{i, j\}$, we may assume

$$\tilde{\alpha}_{\{i,j\}}^d = \tau_{i,j}^d = d(i, j) + d(k, l) - d(i, k) - d(j, l). \quad (4.38)$$

From the conditions (C2) and (C3), the possible partial splits in \mathcal{S} that contribute to the term $d(i, k) + d(j, l)$ are five partial splits $\{i, jkl\}$, $\{j, ikl\}$, $\{i, jk\}$, $\{j, il\}$ and $\{il, jk\}$. But these partial split distances cancel out in (4.38). Therefore we obtain $\tilde{\alpha}_{\{i,j\}}^d \geq \lambda_{\{i,j\}}$. Hence we conclude that (4.35) holds for $|V| = 4$. \square

By Theorem 4.13 and Proposition 3.5, weak compatibility of a collection of splits is also characterized by Λ -admissibility of the corresponding collection of hyperplanes.

Theorem 4.14. *For a collection of partial splits \mathcal{S} , the following conditions (a) and (b) are equivalent.*

- (a) \mathcal{S} is weakly compatible.
- (b) $\{H_S \mid S \in \mathcal{S}\}$ is Λ -admissible.

Figure 11 illustrates all types of maximal weakly compatible partial splits on 4-point set and corresponding hyperplane arrangements.

Remark 4.15. The fundamental fact [1, Corollary 4] that the number of weakly compatible splits is bounded by $|V|(|V| - 1)/2$ also follows from Proposition 3.6. It is shown that this bound is attained by the maximum *circular split system*, which is obtained from a convex $|V|$ -gon [1, Theorem 5]. By Corollary 3.7, the sum of maximum circular split metrics yields a triangulation of $\text{conv } \Lambda$. We point out that this construction of a triangulation of $\text{conv } \Lambda$ is essentially equivalent to the construction of the triangulation of the second hypersimplex $\text{conv}\{\chi_i + \chi_j \mid i, j \in V, i \neq j\}$ due to De Loera, Sturmfels, and Thomas [6] (see also [25, Chapter 9]).

A collection of splits \mathcal{S} is said to be *compatible* if for any pair of splits $\{A, B\}, \{C, D\} \in \mathcal{S}$, at least one of four sets $A \cap C, A \cap D, B \cap C$ and $B \cap D$ is empty (see [5],[2],[24]). Compatibility of a collection of splits can also be captured as a geometric property of the corresponding collection of hyperplanes.

Proposition 4.16. *For a collection of splits \mathcal{S} , the following conditions (a) and (b) are equivalent.*

(a) \mathcal{S} is compatible.

(b) $H_1 \cap H_2 \cap \text{ri conv } \Lambda = \emptyset$ holds for each pair $H_1, H_2 \in \{H_S \mid S \in \mathcal{S}\}$.

Proof. For two splits $\{A, B\}, \{C, D\} \in \mathcal{S}$, consider two hyperplanes $H_{\{A,B\}}$ and $H_{\{C,D\}}$. The nonemptiness of $H_{\{A,B\}} \cap H_{\{C,D\}} \cap \text{ri conv } \Lambda$ is equivalent to the existence of a solution to linear equality, inequality system

$$\begin{cases} x(A) - x(B) = 0, \\ x(C) - x(D) = 0, \\ x(i) > 0 \quad (i \in V). \end{cases} \quad (4.39)$$

(a) \Rightarrow (b). By compatibility of splits \mathcal{S} , we may assume $A \subset C$ and $D \subset B$. By subtracting the second of (4.39) from the first, we have $x(B \cap C) = 0$ and hence (4.39) is empty. (b) \Rightarrow (a). Suppose that $\{A, B\}$ and $\{C, D\}$ are incompatible, i.e., all the four sets $A \cap C, A \cap D, B \cap C$ and $B \cap D$ are nonempty. Then $x \in \mathbf{R}^V$ defined by

$$x(i) = \begin{cases} 1/|A \cap C| & \text{if } i \in A \cap C \\ 1/|A \cap D| & \text{if } i \in A \cap D \\ 1/|B \cap C| & \text{if } i \in B \cap C \\ 1/|B \cap D| & \text{if } i \in B \cap D \end{cases} \quad (4.40)$$

is a solution to (4.39). □

A metric d is a *tree metric* if it is represented as the path metric of some weighted tree. It is well known that d is a tree metric if and only if it is represented as

$$d = \sum_{S \in \mathcal{S}} \alpha_S \delta_S \quad (4.41)$$

for some compatible collection of splits \mathcal{S} and a positive weight $\alpha \in \mathbf{R}_{++}^{\mathcal{S}}$ (see [5],[2],[24]). From Propositions 2.9 and 4.16, one of the central theorems in T-theory can be derived.

Theorem 4.17 (Dress [7]). *A metric d is a tree metric if and only if $T(d)$ is a tree.*

Remark 4.18. It is natural to ask when $T(d)$ is a tree for a distance d ; Theorem 4.17 is the case that d is a metric. This question was answered by [15], which shows that the following conditions for a distance $d : V \times V \rightarrow \mathbf{R}$ are equivalent;

- (1) $T(d)$ is a tree.
- (2) d can be represented as $d = \sum_{S \in \mathcal{S}} \alpha_S \zeta_S$ for a compatible collection of partial splits \mathcal{S} and a positive weight $\alpha \in \mathbf{R}_{++}^{\mathcal{S}}$, where a collection of partial splits \mathcal{S} is said to be *compatible* if each pair $\{A, B\}, \{C, D\} \in \mathcal{S}$ satisfies one of the following four conditions:

$$\begin{aligned}
A \subseteq C \text{ and } B \supseteq D, \\
A \subseteq D \text{ and } B \supseteq C, \\
A \supseteq C \text{ and } B \subseteq D, \\
A \supseteq D \text{ and } B \subseteq C.
\end{aligned} \tag{4.42}$$

- (3) There exist some weighted tree T and some collection of subtrees $\{T_i \mid i \in V\}$ of T such that

$$d(i, j) = \text{the shortest path length between } T_i \text{ and } T_j \text{ in } T \quad (i, j \in V) \tag{4.43}$$

(see Figure 3 in the introduction).

Proposition 4.16 can also be extended for a compatible collection of partial splits. In particular, (e) and (f) in Figure 11 are such collections of partial splits. Clearly, a compatible partial splits is weakly compatible. So a distance between subtrees of a tree is totally split-decomposable in our sense and therefore can be recognized by the extended split decomposition.

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