# An algorithm which generates linear extensions for a non-simply-laced d-complete poset with uniform probability 

Kento NAKADA<br>K. NAKADA : Okayama University, Graduate School of Education Master's Program. nakada@okayama-u.ac.jp


#### Abstract

The purpose of this paper is to present an algorithm which generates linear extensions for a non-simplylaced d-complete poset with uniform probability. Résumé. Le but de ce papier est présenter un algorithme qui produit des extensions linéaires pour une non-simplylaced d-complete poset avec probabilité constante.


Keywords: d-complete posets, algorithm, linear extension, uniform generation

## 1 Introduction

In [7](Theorem 4.2), J. Stembridge classified irreducible minuscule elements of Kac-Moody Weyl group over a root system $\Phi$ into three classes below:

- $\Phi$ is simply-laced,
- $\Phi$ has the form

(namely, of type $B$ ), or
- $\Phi$ has the form

In [5] [6], the author and S. Okamura constracted an algorithm which generates reduced decompositions for a given minuscule element of simply-laced Weyl group with uniform probability. The algorithm in [6] is described in terms of graphs. Simply-laced minuscule elements are described as certain simple acyclic di-graphs. The transitive-closure of the graph is called a d-complete poset. Then, the reduced decompositions are identified with linear extensions of the graph. This algorithm gives a proof of the hook formula [1] for the number of reduced decompositions of a minuscule element in simply-laced case.

In this paper, we present an algorithm (algorithm A) in terms of graphs (See Section 2 for details). This algorithm is a generalization of an algorithm in [5][6]. We define a certain acyclic multi-di-graph corresponding to a minuscule element of type $B$ (resp. type $F_{m}$ ) in Section 3 (resp. Section 4). Our main result (Theorem5.1) is that the algorithm A generates linear extensions for a minuscule element of
type $B$ and $F_{m}$ with uniform probability. More precisely, the probability the algorithm A generates linear extension $L$ of a graph $S$ is given by:

$$
\begin{equation*}
\frac{\prod_{v \in S}\left(1+\# \mathrm{H}_{S}(v)^{+}\right)}{\# S!} \tag{1.1}
\end{equation*}
$$

where $\mathrm{H}_{S}(v)^{+}$is a certain subset of $S$ (See Section 2 for detail). This 1.1) is independent from the choice of $L$. Hence, we get the hook formula for the number of linear extensions of a given shape $S$ of type $B$ and $F_{m}$. Namely, the number of linear extensions of a shape $S$ is given by:

$$
\frac{\# S!}{\prod_{v \in S}\left(1+\# \mathrm{H}_{S}(v)^{+}\right)}
$$

In section 6 , we give a Lie theoretical description of shape of type $B$ and $F_{m}$.

## 2 An algorithm for a graph $\Gamma$

Let $\Gamma=(\Gamma ; A, \mathrm{o}, \mathrm{i})$ be a finite acyclic multi-di-graph, where $A$ denotes the set of arrows of $\Gamma, \mathrm{i}(a)$ the sink of $a \in A$, and o $(a)$ the source of $a \in A$.

Definition 2.1 Put $d:=\# \Gamma$. A bijection $L:\{1, \cdots, d\} \longrightarrow \Gamma$ is said to be a linear extension of $\Gamma$ if:

$$
L(k)=\mathrm{o}(a) \text { and } \mathrm{i}(a)=L(l) \text { implies } k>l, \quad k, l \in\{1, \cdots, d\}, \quad a \in A
$$

The set of linear extensions of $\Gamma$ is denoted by $\mathcal{L}(\Gamma)$.
For a given $v \in \Gamma$, we define a set $\mathrm{H}_{\Gamma}(v)^{+}$by:

$$
\mathrm{H}_{\Gamma}(v)^{+}:=\{a \in A(\Gamma) \mid v=\mathrm{o}(a)\}
$$

For a given $\Gamma$, we call the following algorithm the algorithm A for $\Gamma$ :
GNW1. Set $i:=0$ and set $\Gamma_{0}:=\Gamma$.
GNW2. (Now $\Gamma_{i}$ has $d-i$ nodes.) Set $j:=1$ and pick a node $v_{1} \in \Gamma_{i}$ with the probability $1 /(d-i)$.
GNW3. If $\# \mathrm{H}_{\Gamma_{i}}\left(v_{j}\right)^{+} \neq 0$, pick an arrow $a_{j+1} \in \mathrm{H}_{\Gamma_{i}}\left(v_{j}\right)^{+}$with the probability $1 / \# \mathrm{H}_{\Gamma_{i}}\left(v_{j}\right)^{+}$. If not, go to GNW5.

GNW4. Set $v_{j+1}:=\mathrm{i}\left(a_{j}\right)$. Set $j:=j+1$ and return to GNW3.
GNW5. (Now $\# \mathrm{H}_{\Gamma_{i}}\left(v_{j}\right)^{+}=0$.) Set $L(i+1):=v_{j}$ and set $\Gamma_{i+1}:=\Gamma_{i} \backslash v_{j}$ (the graph deleted $v_{j}$ from $\Gamma_{i}$ ).
GNW6. Set $i:=i+1$. If $i<d$, return to GNW2; if $i=d$, terminate.
We note that the algorithm A stops in finite time since $\Gamma$ is acyclic. By the definition of the algorithm A for $\Gamma$, the map $L: i \mapsto L(i)$ generated above is a linear extension of $\Gamma$. We denote by $\operatorname{Prob}_{\Gamma}(L)$ the probability we get $L \in \mathcal{L}(\Gamma)$ by the algorithm A .

## 3 Shapes of type $B$

We denote by $\mathbb{N}$ the set of non-negative integers. We define a set $\mathbb{B}$ by:

$$
\mathbb{B}:=\{(i, j) \in \mathbb{N} \times \mathbb{N} \mid i \leq j\}
$$

The set $\mathbb{B}$ is depicted in FIGURE 3.1. We equip the $\mathbb{B}$ with the partial order:

$$
(i, j) \leq\left(i^{\prime}, j^{\prime}\right) \Longleftrightarrow i \geq i^{\prime} \text { and } j \geq j^{\prime}
$$


$\ddots$.
Fig. 3.1: The set $\mathbb{B}$

Definition 3.1 Let $S$ be a finite order filter of $\mathbb{B}$. We induce to $S$ a graph structure by:

$$
\begin{aligned}
& (i, j) \rightarrow\left(i^{\prime}, j^{\prime}\right) \quad \text { if and only if }\left\{\begin{array}{c}
" i=j \text { and } i^{\prime}=i, j^{\prime}>j ", \\
" i<j \text { and } i^{\prime}=i, j^{\prime}>j ", \\
" i<j \text { and } i^{\prime}>i, j^{\prime}=j ", \\
\text { or " } i<j \text { and } i^{\prime}=j, j^{\prime}>i ",
\end{array}\right. \\
& (i, j) \rightrightarrows\left(i^{\prime}, j^{\prime}\right) \text { if and only if " } i<j \text { and } i^{\prime}=j^{\prime}=j ",
\end{aligned}
$$

and there exists no other adjacency relation. Here, $v \rightarrow v^{\prime}$ means there exists exactly one arrow from $v$ to $v^{\prime}$, and $v \rightrightarrows v^{\prime}$ there exists exactly two arrows from $v$ to $v^{\prime}$. A graph $S$ is called a shape of type B. See FIGURE 3.2 for examples of $\mathrm{H}_{S}(v)^{+}$.


Fig. 3.2: $\mathrm{H}_{S}(u)^{+}, \mathrm{H}_{S}(v)^{+}$, and $\mathrm{H}_{S}(w)^{+}$.

Remark 3.2 A shape of type B as poset is order-isomorphic to a shifted shape. Shifted shapes are also realized as d-complete posets over a root system of type $D$. The graph-structure of shapes of type $D$ is described in [6] and compatible with notion of hooks (or called bars) of shifted shapes. The algorithm A depends not only on poset-structure but on graph-structure. Hence, we do not consider shapes of type $B$ as shifted shapes.

4 Shapes of type $F_{m}(m \geq 2)$.
We denote by $\mathbb{Z}$ the set of integers. Let $m$ be an integer greater than or equal to 2 . We define a set $\mathbb{F}_{m}$ by:

$$
\mathbb{F}_{m}:=\left\{\begin{array}{l|l}
(i, j) \in \mathbb{N} \times \mathbb{Z} & \begin{array}{l}
i=0 \text { and } j \geq-m \\
i=1 \text { and } j \geq 0, \text { or } \\
2 \leq i \leq m \text { and } j=0
\end{array}
\end{array}\right\}
$$

For example, the set $\mathbb{F}_{3}$ is depicted in FIGURE 4.1. We equip the $\mathbb{F}_{m}$ with the partial order:

$$
(i, j) \leq\left(i^{\prime}, j^{\prime}\right) \Longleftrightarrow i \geq i^{\prime} \text { and } j \geq j^{\prime}
$$



Fig. 4.1: The set $\mathbb{F}_{3}$

Definition 4.1 Let $S$ be a finite order filter of $\mathbb{F}_{m}$. We induce to $S$ a graph structure by:

$$
\begin{aligned}
& (i, j) \rightarrow\left(i^{\prime}, j^{\prime}\right) \text { if and only if }\left\{\begin{array}{l}
" i=0, j \leq-1 \text { and } i^{\prime} \neq-j, j^{\prime}>j ", \\
" i=0, j=0, \text { and } j^{\prime}>0 ", \\
" i=1, j=0, \text { and } i^{\prime}=1, j^{\prime}>0 ", \\
" i=1, j=0, \text { and } i^{\prime}>1, j^{\prime}=0 ", \\
" i \geq 2, j=0, \text { and } i^{\prime}>i, j^{\prime}=0 ", \\
" j \geq 1 \text { and } i^{\prime}=i, j^{\prime}>j^{\prime} " \\
\text { or " } j \geq 1 \text { and } i^{\prime}>i, j^{\prime}=j ",
\end{array}\right. \\
& (i, j) \rightrightarrows\left(i^{\prime}, j^{\prime}\right) \text { if and only if } " i=0, j=0, \text { and } 0<i^{\prime}, j^{\prime}=0 ",
\end{aligned}
$$

and there exists no other adjacency relation. A graph $S$ is called a shape of type $F_{m}$. See FIGURE 4.2 for examples of $\mathrm{H}_{S}(v)^{+}$.


Fig. 4.2: $\mathrm{H}_{S}(u)^{+}, \mathrm{H}_{S}(v)^{+}$, and $\mathrm{H}_{S}(w)^{+}$.

## 5 Main result

Now, we can state the main theorem:
Theorem 5.1 Let $S$ be a shape of type $B$ or type $F_{m}$ for some $m \geq 2$. Let $L \in \mathcal{L}(S)$. Then the algorithm A for $S$ generates $L$ with the probability

$$
\begin{equation*}
\operatorname{Prob}_{S}(L)=\frac{\prod_{v \in S}\left(1+\# \mathrm{H}_{S}(v)^{+}\right)}{\# S!} \tag{5.1}
\end{equation*}
$$

Since the right hand side of (5.1) is independent from the choice of $L \in \mathcal{L}(S)$, we have:
Corollary 5.2 Let $S$ be a shape of type $B$ or type $F_{m}$ for some $m \geq 2$. Then we have:

$$
\# \mathcal{L}(S)=\frac{\# S!}{\prod_{v \in S}\left(1+\# \mathrm{H}_{S}(v)^{+}\right)}
$$

## 6 Lie theoretical description of main result and Remarks

In this section, we fix a (not necessary simply-laced) Kac-Moody Lie algebra $\mathfrak{g}$ with a simple root system $\Pi=\left\{\alpha_{i} \mid \in I\right\}$. For all undefined terminology in this section, we refer the reader to [2] [3].
Definition 6.1 An integral weight $\lambda$ is said to be pre-dominant if:

$$
\left\langle\lambda, \beta^{\vee}\right\rangle \geq-1 \quad \text { for each } \beta^{\vee} \in \Phi_{+}^{\vee},
$$

where $\Phi_{+}^{\vee}$ denotes the set of positive real coroots. The set of pre-dominant integral weights is denoted by $P_{\geq-1}$. For $\lambda \in P_{\geq-1}$, we define the set $\mathrm{D}(\lambda)^{\vee}$ by:

$$
\mathrm{D}(\lambda)^{\vee}:=\left\{\beta^{\vee} \in \Phi_{+}^{\vee} \mid\left\langle\lambda, \beta^{\vee}\right\rangle=-1\right\}
$$

The set $\mathrm{D}(\lambda)^{\vee}$ is called the shape of $\lambda$. If $\# \mathrm{D}(\lambda)^{\vee}<\infty$, then $\lambda$ is called finite.
Proposition 6.2 (see [4]) Let $\lambda \in P_{\geq-1}$ be finite and $\beta^{\vee}, \gamma^{\vee} \in D(\lambda)^{\vee}$ satisfy $\beta^{\vee}>\gamma^{\vee}$ in the ordinary order of coroots. Then we have:

$$
\left\langle\beta, \gamma^{\vee}\right\rangle=0,1, \text { or } 2
$$

By proposition 6.2, we introduce graph-structure into $\mathrm{D}(\lambda)^{\vee}$ by:

$$
\begin{aligned}
& \beta^{\vee} \rightarrow \gamma^{\vee} \Leftrightarrow \beta^{\vee}>\gamma^{\vee} \text { and }\left\langle\beta, \gamma^{\vee}\right\rangle=1 \\
& \beta^{\vee} \rightrightarrows \gamma^{\vee} \Leftrightarrow \beta^{\vee}>\gamma^{\vee} \text { and }\left\langle\beta, \gamma^{\vee}\right\rangle=2
\end{aligned}
$$

If $\beta^{\vee} \ngtr \gamma^{\vee}$, or $\beta^{\vee}>\gamma^{\vee}$ and $\left\langle\beta, \gamma^{\vee}\right\rangle=0$, then no arrows from $\beta^{\vee}$ to $\gamma^{\vee}$ exist.
Thus, we get a finite acyclic multi-di-graph $\mathrm{D}(\lambda)^{\vee}$ for a finite $\lambda \in P_{\geq-1}$.
Remark 6.3 The finite pre-dominant integral weights $\lambda$ are identified with the minuscule elements $w[4]$. And, we have $\mathrm{D}(\lambda)^{\vee}=\left\{\beta^{\vee} \in \Phi_{+}^{\vee} \mid w^{-1}\left(\beta^{\vee}\right)<0\right\}$. Furthermore, the linear extensions of $\mathrm{D}(\lambda)^{\vee}$ are identified with the reduced decompositions of $w[4]$ by the following one-to-one correspondence:
$\operatorname{Red}(w) \ni\left(s_{i_{1}}, s_{i_{2}}, \cdots, s_{i_{d}}\right) \longleftrightarrow L \in \mathcal{L}\left(\mathrm{D}(\lambda)^{\vee}\right), \quad L(k)=s_{i_{1}} \cdots s_{i_{k-1}}\left(\alpha_{i_{k}}\right)^{\vee} \in \mathrm{D}(\lambda)^{\vee}(k=1, \cdots d)$,
where $\operatorname{Red}(w)$ denotes the set of reduced decompositions of $w, d=\ell(w)$ the length of $w$.

### 6.1 Case of type $B$

Suppose that the underlying Dynkin diagram is of type $B$ :


Let $W=\left\langle s_{0}, s_{1}, s_{2}, \cdots\right\rangle$ be the Weyl group. Let $\Lambda_{0}$ be the 0 -th fundamental weight. Then each $\lambda \in W \Lambda_{0}$ is a finite pre-dominant integral weight. And, $\mathrm{D}(\lambda)^{\vee}$ is graph-isomorphic with some shape of type $B$ defined in section 3 .

Remark 6.4 Let $W_{0}:=\left\langle s_{1}, s_{2}, \cdots\right\rangle$ be a maximal parabolic subgroup of $W$, which is the Weyl group of type $A$. Then a minimal coset representative $w$ in $W / W_{0}$ is called a Lagrangian Grassmannian element.
Let $\lambda \in W \Lambda_{0}$. Then the corresponding minuscule element $w$ in remark 6.3 is a Lagrangian Grassmannian element. Our result gives the number of reduced decompositions of Lagrangian Grassmannian element $w$.

### 6.2 Case of type $F_{m}(m \geq 2)$

Let $m \in \mathbb{Z}$ be greater than or equal to 2 . Suppose that the underlying Dynkin diagram is of type $F_{m}$ :


Let $W=\left\langle s_{-m}, \cdots, s_{-2}, s_{-1}, s_{0}, s_{1}, \cdots\right\rangle$ be the Weyl group. Let $\Lambda_{-m}$ be the $(-m)$-th fundamental weight. Then each $\lambda \in P_{\geq-1} \cap W \Lambda_{-m}$ is a finite pre-dominant integral weight. And, $\mathrm{D}(\lambda)^{\vee}$ is graphisomorphic with some shape of type $F_{m}$ defined in section 4.

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